

CONSERVATION AS TRANSLATION[†]

GIULIO FELLIN 

Dipartimento di Informatica, Università degli Studi di Verona
Dipartimento di Ingegneria dell'Informazione, Università degli Studi di Brescia
and

PETER SCHUSTER 

Dipartimento di Informatica, Università degli Studi di Verona

Abstract. Glivenko's theorem says that classical provability of a propositional formula entails intuitionistic provability of the double negation of that formula. This stood right at the beginning of the success story of negative translations, indeed mainly designed for converting classically derivable formulae into intuitionistically derivable ones. We now generalise this approach: simultaneously from double negation to an arbitrary nucleus; from provability in a calculus to an inductively generated abstract consequence relation; and from propositional logic to any set of objects whatsoever. In particular, we give sharp criteria for the generalisation of classical logic to be a conservative extension of the one of intuitionistic logic with double negation.

§1. Introduction. Glivenko's theorem says that, in propositional logic, classical provability of a formula entails intuitionistic provability of the double negation of that formula [45]. This stood right at the beginning of the success story of negative translations, indeed mainly designed for converting classically derivable formulae into intuitionistically derivable ones. Typically, a negative translation is an operation on formulae in predicate logic such that each formula is classically equivalent to its translation and is classically derivable exactly when the translation is intuitionistically derivable. Negative translations have already been put into the context of nuclei [30, 114] or nucleus [30], and have proved to be useful also in computer science [48], set theory [4], arithmetic and analysis [107], eventually contributing to program extraction [101] and proof mining [63]. Double negation over intuitionistic logic is indeed a typical instance of a nucleus [5, 50, 60, 76, 93, 105, 106, 114].¹

As compared to recent literature on related topics [4, 15, 31, 36, 40, 48, 49, 57, 68, 77, 78, 83, 85, 94, 107, 111, 119], the main purpose of the present work is to

Received: June 19, 2023.

2020 *Mathematics Subject Classification*: Primary 03F50, Secondary 03B55, 03G27.

Key words and phrases: Glivenko, nucleus, entailment relation, negative translation, double negation.

[†]This note grew out from a chapter of the first author's doctoral thesis [33], and is a revised and extended version of a conference paper [35]. As compared to the latter, the main amendments include: the study of j -homogeneous functions and j -translations; a more informative statement of the Gödel conservation theorem; the new Kolmogorov conservation theorem and its corollary, the Kuroda conservation theorem; and an application to lax logic.

¹ This and other lists of references below are meant rather indicative than exhaustive.

© The Author(s), 2025. Published by Cambridge University Press on behalf of The Association for Symbolic Logic. This is an Open Access article, distributed under the terms of the Creative Commons Attribution licence (<https://creativecommons.org/licenses/by/4.0/>), which permits unrestricted re-use, distribution, and reproduction in any medium, provided the original work is properly cited.



generalise Glivenko’s theorem and the theory of negative translations: simultaneously from double negation to an arbitrary nucleus; from provability in a calculus to an abstract consequence relation; and from propositional logic to any set of objects whatsoever. As pointed out by van den Berg [114] and Escardó–Oliva [29, 30], a generalised version of the Gentzen negative translation for arbitrary nuclei is already known in logic [4], locale theory [60] and topos theory [62]; and van den Berg himself gives a generalisation of the minimal Kuroda negative translation to nuclei in logic. We aim for a deeper insight into these generalisations, dubbed j -translations, by passing to arbitrary sets endowed with abstract consequence relations.²

To this end we move to a nucleus j over a Hertz–Tarski consequence relation in the form of a (single-conclusion) entailment relation \triangleright à la Scott [14, 103]. Assuming that \triangleright is inductively generated by axioms and rules, we propose two natural extensions (§2.2): \triangleright_j generalises the provability of double negation, and \triangleright^j is inductively defined by adding the generalisation of double negation elimination to the inductive definition of \triangleright . By their very definitions, \triangleright^j satisfies all the axioms and rules of \triangleright , and \triangleright_j satisfies all the axioms of \triangleright . But when does \triangleright_j also satisfy all the rules of \triangleright ? Glivenko conservation, Corollary 3.7, says that \triangleright^j extends \triangleright_j , and that the two relations coincide precisely when \triangleright_j is closed under the non-axiom rules used to inductively generate \triangleright , which of course is the case whenever there are no such non-axiom rules (Corollary 3.8). While it may happen that this closure condition fails, and hence \triangleright^j is not conservative over \triangleright_j , following Gödel one can ensure conservation by adding to \triangleright suitable rules generalising the double negation shift (Theorem 3.5).

We then investigate into a different generalisation by weakening the conditions on the function: instead of a nucleus j , we consider a j -homogeneous function k and the related entailment relation $\triangleright_{(k)}$ which generalise the Kolmogorov negative translation and its provability. Similarly to the above, we notice that \triangleright^j extends $\triangleright_{(k)}$, and that the two relations coincide precisely when $\triangleright_{(k)}$ is closed under the rules that are used to inductively generate \triangleright (Theorem 3.13). A variant of this last result is shown for $k = jJ$ —which is intended to generalise the Kuroda negative translation—where J is j -homogeneous (Corollary 3.14).

In logic, the prime instance of course is Glivenko’s theorem (Proposition 5.4) as a syntactical conservation theorem (see also [36]):

$$\Gamma \triangleright_c \varphi \iff \Gamma \triangleright_i \neg\neg\varphi,$$

where \triangleright_c and \triangleright_i denote classical and intuitionistic propositional logic. Simultaneously we re-obtain Gödel’s theorem [46], otherwise ascribed to Gabbay [39] (Proposition 5.7), which says that, in predicate logic,

$$\Gamma \triangleright_c \varphi \iff \Gamma \triangleright_* \neg\neg\varphi,$$

where \triangleright_* is any extension (by additional axioms) of intuitionistic predicate logic that satisfies the double negation shift:

$$\forall x \neg\neg\varphi \triangleright \neg\neg\forall x\varphi.$$

² Most recently van den Berg [115] has put forward a theory of nuclei for Miquel’s implicative algebras [72].

Next, we regain Kolmogorov’s result [64] (Proposition 5.8(i)) that, in predicate logic,

$$\Gamma \triangleright_c \varphi \iff k\Gamma \triangleright_i k\varphi,$$

where k is the Kolmogorov negative translation inductively defined by

$$\begin{aligned} k\top &= \neg\neg\top, & k\perp &= \neg\neg\perp, & kP &= \neg\neg P, \\ k(\varphi \wedge \psi) &= \neg\neg(k\varphi \wedge k\psi), & k(\varphi \vee \psi) &= \neg\neg(k\varphi \vee k\psi), & k(\varphi \rightarrow \psi) &= \neg\neg(k\varphi \rightarrow k\psi), \\ k(\forall x \varphi) &= \neg\neg\forall x k\varphi, & k(\exists x \varphi) &= \neg\neg\exists x k\varphi. \end{aligned}$$

Analogous statements can be made for the Gentzen [41, 42] and for the minimal Kuroda [66, 73, 114] negative translation (Proposition 5.8(ii) and (iii)).

While the double negation nucleus $j\varphi = \neg\neg\varphi$ is an instance of the continuation nucleus, it is tantamount to the same case $j\varphi = \neg\varphi \rightarrow \varphi$ of the Peirce nucleus [30]. What does our main result mean for other nuclei in logic? The open nucleus $j\varphi = A \rightarrow \varphi$ prompts a form of the deduction theorem (Proposition 5.11(i)), while the closed nucleus $j\varphi = \varphi \vee A$ yields a variant of the reduction from intuitionistic to minimal logic going back to Johansson [59] (Proposition 5.11(ii)). Last but not least, by observing that the modal operator \bigcirc in propositional lax logic is a nucleus, we get a general version of strong conservativity [32] (Proposition 5.13).

1.1. Preliminaries. We proceed in a constructive and predicative way, keeping the concepts elementary and the proofs direct. If a formal system is desired, our work can be placed in a suitable fragment of Aczel’s *Constructive Zermelo–Fraenkel Set Theory (CZF)* [1–3, 6, 7] based on intuitionistic first-order predicate logic.

By a *finite set* we understand a set that can be written as $\{a_1, \dots, a_n\}$ for some $n \geq 0$. Given any set S , let $\text{Pow}(S)$ (respectively, $\text{Fin}(S)$) consist of the (finite) subsets of S . We thus deviate from the terminology prevalent in constructive mathematics and set theory [6, 7, 11–13, 69, 71]: to reserve the term ‘finite’ to sets which are in *bijection* with $\{1, \dots, n\}$ for a necessarily unique $n \geq 0$. Those exactly are the sets which are finite in our sense and are *discrete* too, i.e., have decidable equality [71].

§2. Entailment relations. Closely following [35, 89, 90] we briefly recall the basics of entailment relations. Let S be a set and $\triangleright \subseteq \text{Pow}(S) \times S$. Once abstracted from the context of logical formulae, all but one of Tarski’s axioms of *consequence* [109]³ can be put as

$$\frac{U \ni a}{U \triangleright a} \quad \frac{\forall b \in U (V \triangleright b) \quad U \triangleright a}{V \triangleright a} \quad \frac{U \triangleright a}{\exists U_0 \in \text{Fin}(U) (U_0 \triangleright a)},$$

where $U, V \subseteq S$ and $a \in S$. These axioms also characterise a finitary covering or Stone covering in formal topology [95];⁴ see further [17, 18, 75, 76, 96, 97]. The notion of consequence has presumably been described first by Hertz [51–53] (see also [9, 67]).

³ Tarski has further required that S be countable.

⁴ This is from where we have taken the symbol \triangleright , used also [16, 116] to denote a ‘consecution’ [87].

Tarski has rather characterised the set of consequences of a set of propositions, which corresponds to the *algebraic closure operator* $U \mapsto U^\triangleright$ on $\text{Pow}(S)$ of a relation \triangleright as above where

$$U^\triangleright = \{a \in S : U \triangleright a\}.$$

Rather than with Tarski’s notion, we henceforth work with its (tantamount) restriction to finite subsets, i.e., a (*single-conclusion*) *entailment relation*.⁵ This is a relation $\triangleright \subseteq \text{Fin}(S) \times S$ such that

$$\frac{}{U, a \triangleright a} \text{Refl} \qquad \frac{V \triangleright b \quad V', b \triangleright a}{V, V' \triangleright a} \text{Trans} \qquad \frac{U \triangleright a}{U, U' \triangleright a} \text{Mono}$$

for all finite $U, U', V, V' \subseteq S$ and $a, b \in S$, where as usual $U, V = U \cup V$ and $V, b = V \cup \{b\}$. Our focus thus is on *finite* subsets of S , for which we reserve the letters U, V, W, \dots ; we sometimes write a_1, \dots, a_n in place of $\{a_1, \dots, a_n\}$ even if $n = 0$.

REMARK 2.1. *The rule Refl is equivalent, by Mono, to the axiom $a \triangleright a$.*

Redefining

$$T^\triangleright = \{a \in S : \exists U \in \text{Fin}(T)(U \triangleright a)\}$$

for *arbitrary* subsets T of S gives back an algebraic closure operator on $\text{Pow}(S)$. By writing $T \triangleright a$ in place of $a \in T^\triangleright$, the entailment relations thus correspond exactly to the relations satisfying Tarski’s axioms above.

Given an entailment relation \triangleright , by setting $a \leq b = a \triangleright b$ we get a preorder on S ; whence the conjunction $a \approx b$ of $a \leq b$ and $b \leq a$ is an equivalence relation.

Quite often an entailment relation \triangleright is inductively generated from axioms by closing up with respect to the three *structural rules* Refl, Mono and Trans above [92]. In the present note some leeway is required by allowing for rules other than Refl, Mono and Trans in the inductive generation of \triangleright . These *non-structural rules* have the form⁶

$$\frac{U_1 \triangleright b_1 \quad \dots \quad U_n \triangleright b_n}{U \triangleright b} r \qquad \text{or, more compactly,} \qquad \frac{\{U_i \triangleright b_i : i \leq n\}}{U \triangleright b} r.$$

As an axiom is nothing but a rule with no premiss, we explicitly use *non-axiom rule* to indicate a rule that has at least one premiss, whereas in general we do not excluded that a *rule* be an axiom. If the three structural rules are the only non-axiom rules employed for inductively generating an entailment relation \triangleright , we stress this by saying that \triangleright is *generated only by axioms*.

Given an inductively generated entailment relation \triangleright and a set of rules $*$, then we call \triangleright *plus* $*$ the entailment relation \triangleright_* inductively generated by all the rules that either are used for generating \triangleright or belong to $*$. We also say that \triangleright_* is an *inductive extension* of \triangleright , and we call $*$ the set of *additional rules*.

⁵ In the present note there is no need for abstract *multi-conclusion* consequence or entailment à la Scott [102–104], Lorenzen’s contributions to which are currently under scrutiny [24, 25]. The relevance of multi-conclusion entailment to constructive algebra, point-free topology, etc. has been pointed out in [14], and has widely been used, e.g., in [19–23, 26, 27, 69, 80, 88–91, 98, 99, 117, 118].

⁶ Examples are logical, mathematical and modal rules.

A main feature of inductive generation is that if \triangleright is an entailment relation generated inductively by certain rules, then $\triangleright \subseteq \triangleright'$ for every entailment relation \triangleright' satisfying those rules. By an *extension* \triangleright' of an entailment relation \triangleright we mean in general an entailment relation \triangleright' such that $\triangleright \subseteq \triangleright'$. We say that an extension \triangleright' of \triangleright is *conservative* if also $\triangleright \supseteq \triangleright'$ and thus $\triangleright = \triangleright'$ altogether [36, 89, 90].

2.1. Nuclei over entailment relations. In the context of an entailment relation, what do we mean by a nucleus?

DEFINITION 2.2. *Given a set S endowed with an entailment relation \triangleright , we say that a function $j : S \rightarrow S$ is a nucleus (over \triangleright) if for all $a, b \in S$ and $U \in \text{Fin}(S)$ the following hold:*

$$\frac{U, a \triangleright jb}{U, ja \triangleright jb} Lj \qquad \frac{U \triangleright b}{U \triangleright jb} Rj .$$

REMARK 2.3.

(i) *By Refl and Trans, the rule Rj can be expressed by an axiom, viz.*

$$b \triangleright jb. \tag{1}$$

(ii) *By Lj , from $ja \triangleright ja$ we get $j^2a \triangleright ja$ and thus $j^2a \approx ja$.*

(iii) *The rule Rj is tantamount to the inverse of Lj . In particular, we have*

$$\frac{U \triangleright jb}{jU \triangleright jb} ,$$

where $jU = \{ju : u \in U\}$ and the double line stands for equivalence. This characterisation is already known in particular cases, e.g., the following (iv).

(iv) *If we consider a structure with a relative pseudo-complement operator, i.e., a binary function \rightarrow that satisfies*

$$\frac{U, a \triangleright b}{U \triangleright a \rightarrow b} RPC ,$$

then it is easy to show that nuclei are characterised by a single equivalence:

$$a \rightarrow jb \approx ja \rightarrow jb. \tag{2}$$

This was observed, e.g., in [61, 62, 70] for nuclei in locale theory. In fact, by means of Trans and RPC, (2) is another way of writing Lj and its inverse, which, as noticed in (ii), is tantamount to Rj . Moreover, the following hold by means of RPC, Lj and Rj :

$$j(a \rightarrow b) \triangleright a \rightarrow jb, \tag{3}$$

$$j(a \rightarrow b) \triangleright ja \rightarrow jb. \tag{4}$$

Note that (3) and (4) are equivalent in view of (2). As we will see in the applications (§5), the converses of (3) and (4) do not hold in general.

EXAMPLE 2.4. *The above notion of a nucleus includes as a special case the notion of a nucleus over a locale [5, 58, 60–62, 76, 86, 93, 105, 106], which is well-known as a*

point-free way to put subspaces. In fact, if S is a locale with partial order \leq , then

$$U \triangleright a \iff \bigwedge U \leq a$$

defines an entailment relation [27] such that any given map $j : S \rightarrow S$ is a nucleus over \triangleright precisely when j is a nucleus over the locale S . The latter means that j satisfies

$$ja \wedge jb \leq j(a \wedge b) \tag{5}$$

on top of the conditions for j being a closure operator on S , which can be put as $a \leq ja$ and

$$a \leq jb \implies ja \leq jb. \tag{6}$$

In the presence of $a \leq ja$, which corresponds to Rj in the form of (1), the conjunction of (5) and (6) is equivalent to

$$c \wedge a \leq jb \implies c \wedge ja \leq jb,$$

which in turn subsumes Lj . So the two notions of a nucleus coincide.

EXAMPLE 2.5.

- (i) Every entailment relation \triangleright has the trivial nucleus $j = \text{id}$.
- (ii) Let an algebraic structure \mathbf{S} come with a unary self-inverse function j (e.g., take a group as \mathbf{S} and the inverse as j). The entailment relation \triangleright of \mathbf{S} -substructures is inductively defined by

$$a_1, \dots, a_n \triangleright f(a_1, \dots, a_n) \tag{7}$$

for every n -ary function f characteristic of \mathbf{S} , including j . Then j is a nucleus over \triangleright . Axiom (1) is just (7) for $f = j$, therefore the rule Rj holds. In particular, $j^2 = \text{id}$ implies $j(a) \triangleright a$, which, together with Trans , gives the rule Lj .

- (iii) Double negation $\neg\neg$ is a nucleus over intuitionistic logic \triangleright_i as an entailment relation between propositional or first-order predicate formulae (see §5.2 for further details).

2.2. Extensions induced by nuclei. Every nucleus over \triangleright induces two natural extensions of \triangleright as follows.

DEFINITION 2.6. Let j be a nucleus over an entailment relation \triangleright on a set S . We understand by:

- the weak j -extension (or Kleisli extension) of \triangleright the relation $\triangleright_j \subseteq \text{Fin}(S) \times S$ defined by

$$U \triangleright_j a \iff U \triangleright ja;$$

- the strong j -extension of \triangleright the entailment relation $\triangleright^j \subseteq \text{Fin}(S) \times S$ inductively generated by the rules of \triangleright together with the stability axiom for j :

$$ja \triangleright^j a. \tag{8}$$

In the terminology coined before, \triangleright^j is nothing but \triangleright plus the stability axiom for j .

As we will see later on (Corollary 3.7),

$$\triangleright \subseteq \triangleright_j \subseteq \triangleright^j.$$

LEMMA 2.7. *Let S be a set with an entailment relation \triangleright and let j be a nucleus over \triangleright . Then both \triangleright^j and \triangleright_j are entailment relations that extend \triangleright .*

Proof. The statement for \triangleright^j holds by the very definition of \triangleright^j . As for \triangleright_j : By (1) and Remark 2.1, Refl carries over to \triangleright from \triangleright_j . Mono is inherited from \triangleright , and so is Trans in view of Lj:

$$\frac{\frac{U \triangleright ja \quad \frac{V, a \triangleright jb}{V, ja \triangleright jb} Lj}{U, V \triangleright jb} Trans}{U, V \triangleright jb} .$$

Finally, also $\triangleright \subseteq \triangleright_j$ is a consequence of (1). □

REMARK 2.8.

- (i) *Stability and Rj together give $a \approx^j ja$, where \approx^j is the intersection of \triangleright^j and its converse on singletons.*
- (ii) *By Refl in the form of $a \triangleright a$ (Remark 2.1), stability holds for \triangleright_j too, that is, $ja \triangleright_j a$.*

Under appropriate circumstances, Remark 2.8 will help to obtain $\triangleright^j \subseteq \triangleright_j$ (see Corollaries 3.7 and 3.8).

REMARK 2.9. *The nucleus j on \triangleright is a nucleus also on \triangleright_j and \triangleright^j . In fact, by Lemma 2.7 both extensions inherit axiom (1) from \triangleright , and actually satisfy the following strengthening of Lj:*

$$\frac{U, a \triangleright b}{U, ja \triangleright b} Lj^+ .$$

While Lj^+ for \triangleright_j is just Lj for \triangleright , stability $ja \triangleright a$ is tantamount to Lj^+ for any entailment relation \triangleright whatsoever.

To better understand whether and when \triangleright_j coincides with \triangleright^j , we first study a concrete example.

EXAMPLE 2.10. *Consider deduction in minimal propositional logic \triangleright_m with the closed nucleus $c_\perp : \varphi \mapsto \varphi \vee \perp$ (see §5.3 for details). This \triangleright_m is inductively generated by certain axioms plus the rule*

$$\frac{\Gamma, \varphi \triangleright_m \psi}{\Gamma \triangleright_m \varphi \rightarrow \psi} R\rightarrow$$

which cannot be expressed as an axiom. By its very definition, $\triangleright_m^{c_\perp}$ too satisfies $R\rightarrow$. Does also $\triangleright_{mc_\perp}$ satisfy this rule? If this were the case, then by definition of $\triangleright_{mc_\perp}$ we would have

$$\frac{\Gamma, \varphi \triangleright_m \psi \vee \perp}{\Gamma \triangleright_m (\varphi \rightarrow \psi) \vee \perp} .$$

As $\perp \triangleright_m \psi \vee \perp$, we would obtain $\triangleright_m (\perp \rightarrow \psi) \vee \perp$. However, since minimal logic has the disjunction property and neither disjunct is provable in general, this cannot be the case. So \triangleright_{c_\perp} does not satisfy the rule $R\rightarrow$.

The moral of Example 2.10 is that \triangleright may already have non-axiom rules, such as $R\rightarrow$, which carry over to \triangleright^j by its very definition, and thus need to hold in \triangleright_j too for the former to be conservative over the latter. To deal with this issue, we introduce the following concept.

DEFINITION 2.11. *We say that a rule r that holds for \triangleright is compatible with j over \triangleright if r also holds for \triangleright_j . We usually omit “over j ” if this is clear from the context.*

REMARK 2.12.

- (i) *Refl, Mono, Trans are compatible with every nucleus j , by Lemma 2.7.*
- (ii) *Every composition r of compatible rules is compatible. In fact, the derivation that gives r in \triangleright can be translated smoothly into \triangleright_j , as all the applied rules are compatible.*

This is very useful: to check compatibility for all the rules of an entailment relation \triangleright , it suffices to check compatibility for any set of rules that generate \triangleright .

- (iii) *Each axiom $a_1, \dots, a_n \triangleright b$ can be viewed as a rule with no premiss, and as such is compatible with any given nucleus j , simply by Rj . It follows that, if an entailment relation \triangleright is generated only by axioms, then every rule that holds for \triangleright is compatible with any nucleus j over \triangleright .*

2.3. Homogeneous functions and translations. Throughout this subsection, let j be a nucleus over an entailment relation \triangleright on a set S .

DEFINITION 2.13. *We say that a function $k: S \rightarrow S$ is:*

- *j -homogeneous (from \triangleright^j to \triangleright) if k satisfies the following two conditions:*

$$kja \triangleright ka, \tag{9}$$

$$a \approx^j ka. \tag{10}$$

- *A j -translation (from \triangleright^j to \triangleright) if k satisfies (10) and*

$$\frac{U \triangleright^j b}{kU \triangleright kb}. \tag{11}$$

REMARK 2.14.

- (i) *Condition (10) implies the converse of (11), that is,*

$$\frac{kU \triangleright kb}{U \triangleright^j b}.$$

In fact, since \triangleright^j extends \triangleright , we have that $kU \triangleright kb$ implies $kU \triangleright^j kb$, which is equivalent to $U \triangleright^j b$ by means of $a \approx^j ka$.

- (ii) *The nucleus j is j -homogeneous: it follows immediately from $j^2a \triangleright ja$ (Remark 2.3(ii)) and $a \approx^j ja$ (Remark 2.8).*
- (iii) *Every j -translation is j -homogeneous: by applying (11) to (8) we get (9).*

(iv) *The nucleus j is a j translation precisely when*

$$\frac{U \triangleright^j b}{jU \triangleright jb} .$$

(v) *We have seen that j -translations are j -homogeneous. We want to stress that the two notions are distinct via a counterexample. The closed nucleus $c_{\perp} : \varphi \mapsto \varphi \vee \perp$ over minimal propositional logic \triangleright_m is c_{\perp} -homogeneous by (ii) but it is not a c_{\perp} -translation: this is a direct consequence of Corollary 3.7 and the fact that $R \rightarrow$ is not compatible with c_{\perp} (see Example 2.10).*

(vi) *Every j -translation k that satisfies Rk also satisfies Lk . In fact:*

$$\frac{\frac{a, U \triangleright kb}{a, U \triangleright^j kb} \text{extension} \quad \frac{}{kb \triangleright^j b} (11)}{\frac{a, U \triangleright^j b}{ka, kU \triangleright kb} (10)} \text{Trans} \quad \frac{}{ka, U \triangleright kb} *$$

Where (*) is subsequent applications of Trans with

$$\frac{}{u \triangleright u} \text{Refl} \quad \frac{}{u \triangleright ku} Rk$$

for each $u \in U$. It follows that a j -translation k is a nucleus if and only if it satisfies Rk .

§3. Conservation as translation. We next present a few *conservation theorems for nuclei*, typically giving necessary and sufficient conditions. Throughout this section, let S be a set with an entailment relation \triangleright inductively generated by rules, and let j be a nucleus over \triangleright .

3.1. Glivenko and Gödel conservation.

DEFINITION 3.1. *Given a rule*

$$\frac{\{U_i \triangleright b_i : i \leq n\}}{U \triangleright b} r ,$$

to obtain its j -version r_j we put j in front of every consequent:

$$\frac{\{U_i \triangleright jb_i : i \leq n\}}{U \triangleright jb} r_j .$$

REMARK 3.2.

- (i) *By definition of \triangleright_j , the j -version r_j holds for \triangleright if and only if the original rule r holds for \triangleright_j . This means that if r holds for \triangleright , then r is compatible with j over \triangleright if and only if r_j holds for \triangleright .*
- (ii) *By $j^2a \approx ja$ (Remark 2.3(iii)), r_{jj} holds for \triangleright if and only if r_j holds for \triangleright . In particular, r_j is compatible with j over \triangleright if and only if r_j holds for \triangleright .*
- (iii) *By $ja \approx^j a$ (Remark 2.8), if r holds for \triangleright , then r_j holds for \triangleright^j .*

DEFINITION 3.3. *Given an entailment relation \triangleright generated by rules, let its intermediate j -extension $\triangleright_{\langle j \rangle}$ be \triangleright plus the j -version r_j of all the rules r in the inductive definition of \triangleright .*

REMARK 3.4.

- (i) *We have $\triangleright \subseteq \triangleright_{\langle j \rangle} \subseteq \triangleright^j$ by Remark 3.2(iii).*
- (ii) *By Remarks 2.12 and 3.2(i), to obtain $\triangleright_{\langle j \rangle}$ it suffices to add the j -version of all the non-axiom rules in the inductive definition of \triangleright that are not already compatible with j .*

THEOREM 3.5 (Gödel conservation). *Let \triangleright_* be \triangleright plus additional rules that hold for \triangleright^j and such that j is a nucleus also over \triangleright_* . Then $\triangleright_*^j = \triangleright^j$; in particular, \triangleright^j extends \triangleright_{*j} , that is, $\triangleright_{*j} \subseteq \triangleright^j$. Moreover, the following are equivalent:*

- (a) *\triangleright^j is conservative over \triangleright_{*j} , that is, $\triangleright^j \subseteq \triangleright_{*j}$;*
- (b) *the nucleus j is a j -translation from \triangleright^j to \triangleright_* ;*
- (c) *\triangleright_{*j} is an inductive extension of \triangleright_* ;*
- (d) *all the non-axiom rules that generate \triangleright_* are compatible with j over \triangleright_* , that is, they hold for \triangleright_{*j} ;*
- (e) *\triangleright_* is an extension of $\triangleright_{\langle j \rangle}$, that is, $\triangleright_{\langle j \rangle} \subseteq \triangleright_*$.*

Proof. First, note that $\triangleright \subseteq \triangleright_* \subseteq \triangleright^j$. Since $\triangleright \subseteq \triangleright_*$, we have $\triangleright^j \subseteq \triangleright_*^j$. On the other hand, as stability and all the rules of \triangleright_* hold for \triangleright^j , we also have $\triangleright_*^j \subseteq \triangleright^j$. Therefore $\triangleright_*^j = \triangleright^j$.

Since $a \approx_*^j ja$ by Remark 2.8(i), (b) is tantamount to

$$\frac{U \triangleright^j b}{jU \triangleright_* jb}$$

which by the inverse of Lj (see Remark 2.3(iii)) is equivalent to (a).

As \triangleright^j is an inductive extension of \triangleright_* , (c) follows from (a); and (c) trivially entails (d).

To deduce (a) from (d), assume (d) and consider one by one the rules that generate \triangleright^j : Refl, Mono, Trans hold for \triangleright_j since \triangleright_j is an entailment relation by Lemma 2.7; stability (8) holds for \triangleright_j by Remark 2.8(i); all the rules that generate \triangleright_* hold for \triangleright_j since they are either compatible with j by assumption (d), or axioms and thus compatible with j by Remark 2.12(iii). As \triangleright^j is the smallest extension of \triangleright_* satisfying all these rules, we obtain (a).

By Remark 3.2(ii) with $\triangleright_{\langle j \rangle}$ in place of \triangleright , (d) and thus (a) hold with $\langle j \rangle$ in place of $*$. So, if (e), then

$$\triangleright^j = \triangleright_{\langle j \rangle j} \subseteq \triangleright_{*j} \subseteq \triangleright^j$$

(see also Figure 1); whence (a).

To obtain (e) from (d), suppose (d) and let r be a rule in the inductive definition of \triangleright . By (d), r holds for \triangleright_{*j} which, by Remark 3.2(i) with \triangleright_* in place of \triangleright means that r_j holds for \triangleright_* . Hence all the rules of $\triangleright_{\langle j \rangle}$ hold for \triangleright_* , and thus (e). \square

REMARK 3.6. *Let \triangleright_* be an inductive extension of \triangleright .*

- (i) *For j to be a nucleus over \triangleright_* too, it is sufficient to check Lj . In fact, Rj is tantamount to (1), which as an axiom is inherited by \triangleright_* from \triangleright by the very definition of extension.*

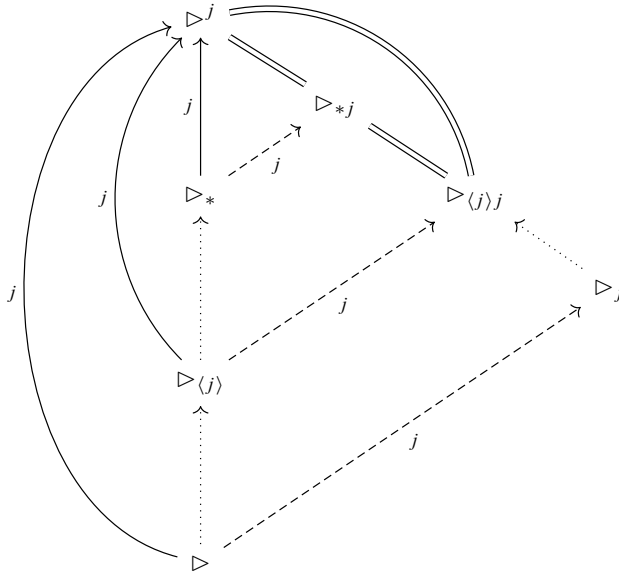


Figure 1. Diagram of the entailment relations involved in the situation of Theorem 3.5. A solid arrow denotes a strong j -extension, a dashed arrow denotes a weak j -extension, a dotted arrow denotes a generic extension, and a double line denotes a conservative extension. The intuition is that, if the outer triangle does not satisfy the desired properties, then we can move to an inner triangle that works.

- (ii) Lj is trivially satisfied whenever $* = \emptyset$.
- (iii) If Lj can be proved admissible for \triangleright by exclusively composing the generating rules of \triangleright , then the same composition equally shows that Lj is admissible for \triangleright_* .

COROLLARY 3.7 (Glivenko conservation). \triangleright^j extends \triangleright_j , that is, $\triangleright_j \subseteq \triangleright^j$; and the following are equivalent:

- (a) \triangleright^j is conservative over \triangleright_j , that is, $\triangleright^j \subseteq \triangleright_j$;
- (b) the nucleus j is a j -translation from \triangleright^j to \triangleright ;
- (c) \triangleright_j is an inductive extension of \triangleright ;
- (d) all the non-axiom rules that generate \triangleright are compatible with j , that is, they hold for \triangleright_j ;
- (e) $\triangleright_{(j)}$ is conservative over \triangleright , that is, $\triangleright_{(j)} \subseteq \triangleright$.

Proof. This is Theorem 3.5 with $\triangleright_* = \triangleright$. □

COROLLARY 3.8. If \triangleright is inductively generated only by axioms, then j is a j -translation and \triangleright^j is a conservative extension of \triangleright_j , that is, $\triangleright^j = \triangleright_j$.

3.2. Kolmogorov and Kuroda conservation. We now generalise Corollary 3.7 from the nucleus j to a j -homogeneous function.

DEFINITION 3.9. For the given nucleus j over an entailment relation \triangleright , and any function $k: S \rightarrow S$ whatsoever, we define the relation $\triangleright_{(k)} \subseteq \text{Fin}(S) \times S$ by

$$U \triangleright_{(k)} a \iff kU \triangleright ka.$$

DEFINITION 3.10. Given a rule r that holds for \triangleright , we say that:

- r is Kolmogorov compatible with k (over \triangleright) if it also holds for $\triangleright_{(k)}$,
- r is j -Kuroda compatible with k (over \triangleright) if it also holds for $\triangleright_{j(k)}$.

We thus have an analogue of Remark 2.12:

REMARK 3.11.

- (i) *Refl, Mono, Trans are trivially Kolmogorov compatible with k . This means that $\triangleright_{(k)}$ is an entailment relation; however in general it is not an extension of the original \triangleright .*
- (ii) *Compatibility entails j -Kuroda compatibility, hence Refl, Mono, Trans are Kuroda compatible with k by Remark 2.12.*
- (iii) *Every composition r of Kolmogorov (resp. j -Kuroda) compatible rules is Kolmogorov (resp. j -Kuroda) compatible. In fact, the derivation that gives r in \triangleright can be translated smoothly into $\triangleright_{(k)}$ (resp. $\triangleright_{j(k)}$), as all the applied rules are Kolmogorov (resp. j -Kuroda) compatible.*

This means that if we want to check Kolmogorov (resp. j -Kuroda) compatibility for all the rules of an entailment relation \triangleright , it suffices to check it for any set of rules that generate \triangleright .

REMARK 3.12. Remark 2.3(iii) yields $\triangleright_{(jk)} = \triangleright_{j(k)}$ and $\triangleright_{(j)} = \triangleright_j$.

THEOREM 3.13 (Kolmogorov conservation). Let $k: S \rightarrow S$ be j -homogeneous. Then \triangleright^j extends $\triangleright_{(k)}$, that is, $\triangleright_{(k)} \subseteq \triangleright^j$; and the following are equivalent:

- (a) \triangleright^j is conservative over $\triangleright_{(k)}$, that is, $\triangleright^j \subseteq \triangleright_{(k)}$;
- (b) the j -homogeneous function k is a j -translation;
- (c) $\triangleright_{(k)}$ is an inductive extension of \triangleright ;
- (d) all the rules in the inductive definition of \triangleright are Kolmogorov compatible with k , that is, they hold for $\triangleright_{(k)}$.

Proof. The fact that $\triangleright_{(k)} \subseteq \triangleright^j$ is just Remark 2.14(i). The implications from (a) to (c) and from (c) to (d) hold just as in Theorem 3.5. Suppose that all the non-structural rules in the inductive definition of \triangleright are Kolmogorov compatible with k , and that $U \triangleright^j b$. We show that $kU \triangleright kb$ by induction on the derivation of $U \triangleright^j b$: the cases involving structural rules are trivial; the case of the stability axiom $ja \triangleright^j a$ is tantamount to $kja \triangleright ka$, which holds since k is j -homogeneous; consider the case in which $U \triangleright^j b$ is derived from a non-structural rule r in the inductive definition of \triangleright , i.e.,

$$\frac{\{U_i \triangleright^j b_i : i \leq n\}}{U \triangleright^j b} r,$$

then

$$\frac{\frac{U_i \triangleright^j b_i}{\{kU_i \triangleright kb_i : i \leq n\}} \text{induction hypothesis}}{kU \triangleright kb} r,$$

where r can be applied because of the Kolmogorov compatibility. This proves that (d) implies (b). Finally, the fact that (b) implies (a) is just (11). □

By setting $k = j$ in Theorem 3.13, by Remark 3.12 we have $\triangleright_{(j)} = \triangleright_j$, which in particular means that compatibility and Kolmogorov compatibility are equivalent. Therefore Corollary 3.7 except for condition (e) is a special case of Theorem 3.13.

COROLLARY 3.14 (Kuroda conservation). *Let S be a set with an entailment relation \triangleright inductively generated by rules, and let j be a nucleus over \triangleright . Let $J : S \rightarrow S$ be j -homogeneous. Then \triangleright^j is an extension of $\triangleright_{j(J)}$, that is, $\triangleright_{j(J)} \subseteq \triangleright^j$; and the following are equivalent:*

- (a) \triangleright^j is conservative over $\triangleright_{j(J)}$, that is, $\triangleright^j \subseteq \triangleright_{j(J)}$;
- (b) the function jJ is a j -translation;
- (c) $\triangleright_{j(J)}$ is an inductive extension of \triangleright ;
- (d) all the non-structural rules in the inductive definition of \triangleright are Kuroda j -compatible with J , that is, they hold for $\triangleright_{j(J)}$.

Proof. First, note that if J is j -homogeneous, then so is $k = jJ$. Then the claim follows immediately from Theorem 3.13 and Remark 3.12. □

§4. Logic as entailment. We now want to see how the conservation results that we have just proved apply to logic.

Throughout this section and the following one, if not stated otherwise, the overall assumption is that S is a set of propositional formulae containing \top, \perp and closed under the connectives $\vee, \wedge, \rightarrow$.⁷ By *minimal (propositional) logic* \triangleright_m we mean the fragment of propositional intuitionistic logic without the principle of *ex falso sequitur quodlibet*. More precisely, we define \triangleright_m as the least entailment relation \triangleright that satisfies the *deduction theorem*

$$\frac{\Gamma, \varphi \triangleright \psi}{\Gamma \triangleright \varphi \rightarrow \psi} \text{R}\rightarrow$$

and the following axioms:

$$\begin{array}{lll} \varphi, \psi \triangleright \varphi \wedge \psi, & \varphi \wedge \psi \triangleright \varphi, & \varphi \wedge \psi \triangleright \psi, \\ \varphi \triangleright \varphi \vee \psi, & \psi \triangleright \varphi \vee \psi, & \varphi \vee \psi, \varphi \rightarrow \delta, \psi \rightarrow \delta \triangleright \delta, \\ \varphi, \varphi \rightarrow \psi \triangleright \psi, & & \\ \triangleright \top. & & \end{array}$$

Of course, we understand this as an inductive definition. In this setting, negation \neg is not given as a primitive operator, but it is rather defined by

$$\neg\varphi = \varphi \rightarrow \perp.$$

The above system for minimal logic is equivalent to the **G3**-style calculus in Table 1; meaning that they inductively generate the same entailment relation.

⁷ It is worth noting that, while we explicitly talk about logic, everything in this section can easily be transferred into any setting with logic-like operators, such as lattice theory, locale theory [60], topos theory [62] and such.

Table 1. *Sequent calculus-like rules for minimal propositional logic.*

$\frac{\Gamma, \varphi, \psi \triangleright \delta}{\Gamma, \varphi \wedge \psi \triangleright \delta} L\wedge$	$\frac{\Gamma \triangleright \varphi \quad \Gamma \triangleright \psi}{\Gamma \triangleright \varphi \wedge \psi} R\wedge$	
$\frac{\Gamma, \varphi \triangleright \delta \quad \Gamma, \psi \triangleright \delta}{\Gamma, \varphi \vee \psi \triangleright \delta} LV$	$\frac{\Gamma \triangleright \varphi}{\Gamma \triangleright \varphi \vee \psi} RV_1$	$\frac{\Gamma \triangleright \psi}{\Gamma \triangleright \varphi \vee \psi} RV_2$
$\frac{\Gamma \triangleright \varphi \quad \Gamma, \psi \triangleright \delta}{\Gamma, \varphi \rightarrow \psi \triangleright \delta} L\rightarrow$	$\frac{\Gamma, \varphi \triangleright \psi}{\Gamma \triangleright \varphi \rightarrow \psi} R\rightarrow$	
$\frac{}{\Gamma \triangleright \top} R\top$		

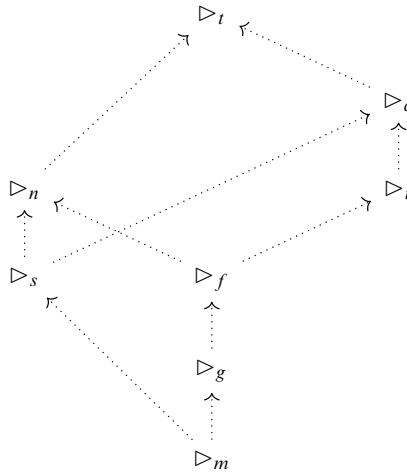


Figure 2. Diagram of the logics introduced in §4, mostly based on [81, figure 4.1].

We define the following inductive extensions of minimal logic \triangleright_m :

extension	additional axiom
<i>intuitionistic logic</i> \triangleright_i	$\perp \triangleright \varphi$ (12)
<i>Glivenko logic</i> \triangleright_g	$\triangleright \neg\neg(\perp \rightarrow \varphi)$ (14)
<i>Clavius logic</i> \triangleright_s	$\neg\varphi \rightarrow \varphi \triangleright \varphi$ (16)
<i>trivial logic</i> \triangleright_t	$\triangleright \varphi$ (18)

extension	additional axiom
<i>classical logic</i> \triangleright_c	$\neg\neg\varphi \triangleright \varphi$ (13)
<i>Frobenius logic</i> \triangleright_f	$\triangleright \perp \vee (\perp \rightarrow \varphi)$ (15)
<i>negative logic</i> \triangleright_n	$\triangleright \perp$ (17)

For a discussion of these and other logics, see [81, 82], from whom we took the names for \triangleright_g and \triangleright_n . See also Figure 2 to see how they are related.

4.1. Predicate and infinitary extensions. Let \triangleright_* be an inductive extension of \triangleright_m . When we say that we work in predicate logic, we mean that we add quantifiers \forall and \exists

Table 2. *Sequent calculus-like rules for quantifiers. Rules $R\forall$ and $L\exists$ come with the condition that y has to be fresh.*

$\frac{\Gamma, \varphi[t/x] \triangleright \delta}{\Gamma, \forall x\varphi \triangleright \delta} L\forall$	$\frac{\Gamma \triangleright \varphi[y/x]}{\Gamma \triangleright \forall x\varphi} R\forall$
$\frac{\Gamma, \varphi[y/x] \triangleright \delta}{\Gamma, \exists x\varphi \triangleright \delta} L\exists$	$\frac{\Gamma \triangleright \varphi[t/x]}{\Gamma \triangleright \exists x\varphi} R\exists$

to the language, and hence require that the formulae in S are closed under \forall, \exists as well. Moreover, we extend the inductive definition of \triangleright_* by adding the rule

$$\frac{\Gamma \triangleright \varphi[y/x]}{\Gamma \triangleright \forall x\varphi} R\forall$$

with the condition that y has to be fresh, and the following axioms:

$$\begin{aligned} &\forall x\varphi \triangleright \varphi[t/x], \\ &\varphi[t/x] \triangleright \exists x\varphi, \\ &\exists x\varphi, \forall x(\varphi \rightarrow \delta) \triangleright \delta, \end{aligned}$$

where the latter comes with the condition that x cannot be free in δ . Over \triangleright_* , these axioms and rule are equivalent to the **G3**-style rules in Table 2.

The definition of a nucleus j given in [114] requires j to be compatible with substitution, that is,

$$j(\varphi[t/x]) = (j\varphi)[t/x].$$

We prefer not to have this as a general assumption, but to make explicit whenever we need it.

When we say that we work in infinitary logic, we mean that we add infinitary connectives $\bigwedge_{i \in \mathbb{N}}$ and $\bigvee_{i \in \mathbb{N}}$ to the language, and hence require that the formulae in S are closed under $\bigwedge_{i \in \mathbb{N}}, \bigvee_{i \in \mathbb{N}}$ as well. Moreover, we extend the inductive definition of \triangleright_* by adding the rule

$$\frac{\{\Gamma \triangleright \varphi_n : n \in \mathbb{N}\}}{\Gamma \triangleright \bigwedge_{i \in \mathbb{N}} \varphi_i} R\wedge$$

and the following axioms:

$$\begin{aligned} &\bigwedge_{i \in \mathbb{N}} \varphi_i \triangleright \varphi_n && \text{for every } n \in \mathbb{N}, \\ &\varphi_n \triangleright \bigvee_{i \in \mathbb{N}} \varphi_i && \text{for every } n \in \mathbb{N}, \\ &\bigvee_{i \in \mathbb{N}} \varphi_i, \bigwedge_{i \in \mathbb{N}} (\varphi_i \rightarrow \delta) \triangleright \delta. \end{aligned}$$

Over \triangleright_* , these axioms and rule are equivalent to the **G3**-style rules in Table 3.

Table 3. *Sequent calculus-like rules for quantifiers. For each $n \in \mathbb{N}$ there is one rule $L\wedge_n$ and one rule $R\vee_n$.*

$\frac{\Gamma, \varphi_n \triangleright \delta}{\Gamma, \bigwedge_{i \in \mathbb{N}} \varphi_i \triangleright \delta} L\wedge_n$	$\frac{\{\Gamma \triangleright \varphi_n : n \in \mathbb{N}\}}{\Gamma \triangleright \bigwedge_{i \in \mathbb{N}} \varphi_i} R\wedge$
$\frac{\{\Gamma, \varphi_n \triangleright \delta : n \in \mathbb{N}\}}{\Gamma, \bigvee_{i \in \mathbb{N}} \varphi_i \triangleright \delta} L\vee$	$\frac{\Gamma \triangleright \varphi_n}{\Gamma \triangleright \bigvee_{i \in \mathbb{N}} \varphi_i} R\vee_n$

§5. Nuclei in logic. Among the logical rules seen in §4, $R\rightarrow$, $R\forall$ and $R\wedge$ are the only ones that cannot be expressed as axioms. Therefore, when checking compatibility of rules with j , if we do not add other rules that cannot be expressed as axioms, then the only rules we have to deal with are these three.

LEMMA 5.1 (Compatibility criterion). *Let \triangleright be an inductive extension of \triangleright_m , and let j be a nucleus over \triangleright .*

- (i) *In either propositional, predicate or infinitary logic, $R\rightarrow$ is compatible with j if and only if*

$$\varphi \rightarrow j\psi \triangleright j(\varphi \rightarrow \psi).$$

- (ii) *In predicate logic, if j is compatible with substitution, that is,*

$$j(\varphi[t/x]) = (j\varphi)[t/x],$$

then $R\forall$ is compatible with j if and only if

$$\forall x j\varphi \triangleright j\forall x\varphi.$$

- (iii) *In infinitary logic, $R\wedge$ is compatible with j if and only if*

$$\bigwedge_{i \in \mathbb{N}} j\varphi_i \triangleright j \bigwedge_{i \in \mathbb{N}} \varphi_i.$$

Proof. We only prove (i), as the proof of (ii) and (iii) is similar. Suppose that $R\rightarrow$ is compatible with j . Then:

$$\frac{\frac{\frac{}{\varphi \rightarrow j\psi, \varphi \triangleright j\psi} \text{axiom}}{\varphi \rightarrow j\psi, \varphi \triangleright_j \psi}}{\varphi \rightarrow j\psi \triangleright_j \varphi \rightarrow \psi} R\rightarrow}{\varphi \rightarrow j\psi \triangleright j(\varphi \rightarrow \psi)} .$$

Now suppose that $\varphi \rightarrow j\psi \triangleright j(\varphi \rightarrow \psi)$. Then:

$$\frac{\frac{\frac{\Gamma, \varphi \triangleright_j \psi}{\Gamma, \varphi \triangleright_j \psi}}{\Gamma \triangleright \varphi \rightarrow j\psi} \text{R}\rightarrow \frac{\Gamma, \varphi \triangleright_j \psi}{\varphi \rightarrow j\psi \triangleright j(\varphi \rightarrow \psi)} \text{assumption}}{\frac{\Gamma \triangleright j(\varphi \rightarrow \psi)}{\Gamma \triangleright_j \varphi \rightarrow \psi} \text{Trans}} \quad \square$$

5.1. Some notable translations in logic. We now consider predicate logic. We introduce some classes of functions that generalise some well-known negative translations. Let \triangleright be either \triangleright_m or \triangleright_i . Given a nucleus j on \triangleright , we inductively define $k, g, t, J : S \rightarrow S$ as follows:

$kP = jP,$	$gP = jP,$	$tP = jP,$	$JP = P,$
$k\top = j\top,$	$g\top = j\top,$	$t\top = j\top,$	$J\top = \top,$
$k\perp = j\perp,$	$g\perp = j\perp,$	$t\perp = j\perp,$	$J\perp = \perp,$
$k(\varphi \wedge \psi) = j(k\varphi \wedge k\psi),$	$g(\varphi \wedge \psi) = g\varphi \wedge g\psi,$	$t(\varphi \wedge \psi) = t\varphi \wedge t\psi,$	$J(\varphi \wedge \psi) = J\varphi \wedge J\psi$
$k(\varphi \vee \psi) = j(k\varphi \vee k\psi),$	$g(\varphi \vee \psi) = j(g\varphi \vee g\psi),$	$t(\varphi \vee \psi) = t\varphi \vee t\psi,$	$J(\varphi \vee \psi) = J\varphi \vee J\psi$
$k(\varphi \rightarrow \psi) = j(k\varphi \rightarrow k\psi),$	$g(\varphi \rightarrow \psi) = g\varphi \rightarrow g\psi,$	$t(\varphi \rightarrow \psi) = t\varphi \rightarrow t\psi,$	$J(\varphi \rightarrow \psi) = J\varphi \rightarrow jJ\psi,$
$k(\forall x \varphi) = j\forall x k\varphi,$	$g(\forall x \varphi) = \forall x g\varphi,$	$t(\forall x \varphi) = \forall x t\varphi,$	$J(\forall x \varphi) = \forall x jJ\varphi,$
$k(\exists x \varphi) = j\exists x k\varphi,$	$g(\exists x \varphi) = j(\exists x g\varphi),$	$t(\exists x \varphi) = \exists x t\varphi,$	$J(\exists x \varphi) = \exists x J\varphi.$

The *Kolmogorov j-function* k is named after the *Kolmogorov negative translation*, which is k obtained for $j = \neg\neg$, as seen in Proposition 5.8(i). A refined version of the Kolmogorov j -function is the *Gentzen j-function* g , named after the *Gentzen negative translation*, which is g obtained for $j = \neg\neg$, as seen in Proposition 5.8(ii). An even more refined version is the *prime j-function* t which, however, does not provide a j -translation for $j = \neg\neg$ as the Gentzen negative translation is known to be minimal in this sense [37]. Finally, we call J the *Kuroda j-function*. Its definition follows van den Berg [114], and is based on the *minimal Kuroda negative translation*, which is jJ for $j = \neg\neg$, as seen in Proposition 5.8(iii).

REMARK 5.2.

- (i) *The following hold for all $\varphi \in S$: $\varphi \approx^j t\varphi$, $\varphi \approx^j k\varphi$, $\varphi \approx^j g\varphi$, $\varphi \approx^j J\varphi$. It follows that t, k, g, J are j -homogenous exactly when they satisfy (9).*
- (ii) *Every rule in the inductive definition of \triangleright is Kolmogorov compatible with t, k, g . This, together with (i) and Theorem 3.13, implies that t, k, g are j -translations exactly when they satisfy (9).*
- (iii) *Every rule in the inductive definition of \triangleright is Kuroda j -compatible with J . This, together with (i) and Corollary 3.14, implies that jJ is a j -translation exactly when J satisfies (9).*

Proof. The properties in (i) are readily seen by induction on φ . We refer to [33, appendix A] for (ii) and (iii).⁸ □

⁸ As both proofs have several cases, they are somewhat lengthy but straightforward otherwise.

5.2. The continuation nucleus and the Glivenko nucleus. Let \triangleright_* be an inductive extension of \triangleright_m , and fix a formula α . The *continuation nucleus*

$$g_\alpha: \varphi \mapsto (\varphi \rightarrow \alpha) \rightarrow \alpha$$

has as special case the *Glivenko nucleus* [93, 114]

$$g_\perp: \varphi \mapsto \neg\neg\varphi.$$

While g_α is well-known to be a nucleus whenever \triangleright_* is \triangleright_m , \triangleright_i or \triangleright_c (see, e.g., [114]), to be sure (Remark 3.6) that this is the case for every inductive extension \triangleright_* of \triangleright_m we give a proof in which we only use generating rules of \triangleright_m ; in fact, Trans and $R\rightarrow$ suffice together with *modus ponens*.

LEMMA 5.3. g_α is a nucleus over \triangleright_* for every formula α .

Proof. Rg_α :

$$\frac{\frac{\frac{\Gamma \triangleright_* \varphi \quad \varphi, \varphi \rightarrow \alpha \triangleright_* \alpha}{\text{Trans}} \text{ axiom}}{\Gamma, \varphi \rightarrow \alpha \triangleright_* \alpha} R\rightarrow}{\Gamma \triangleright_* (\varphi \rightarrow \alpha) \rightarrow \alpha}$$

Lg_α :

$$\frac{\frac{\frac{\Gamma, \varphi \triangleright_* (\psi \rightarrow \alpha) \rightarrow \alpha \quad \psi \rightarrow \alpha, (\psi \rightarrow \alpha) \rightarrow \alpha \triangleright_* \alpha}{\text{Trans}} \text{ axiom}}{\Gamma, \psi \rightarrow \alpha, \varphi \triangleright_* \alpha} R\rightarrow}{\Gamma, \psi \rightarrow \alpha \triangleright_* \varphi \rightarrow \alpha} \frac{\frac{\Gamma, (\varphi \rightarrow \alpha) \rightarrow \alpha, \psi \rightarrow \alpha \triangleright_* \alpha}{\text{Trans}} \text{ axiom}}{\Gamma, (\varphi \rightarrow \alpha) \rightarrow \alpha \triangleright_* (\psi \rightarrow \alpha) \rightarrow \alpha} R\rightarrow \quad \square$$

5.2.1. The intuitionistic case: Glivenko’s theorem. Take propositional intuitionistic logic \triangleright_i as \triangleright . As stability (8) equals double negation elimination (16), the strong extension $\triangleright_i^{g_\perp}$ of intuitionistic logic \triangleright_i is classical logic \triangleright_c .

PROPOSITION 5.4 (Glivenko’s theorem [44, 45]). *In propositional logic,*

$$\Gamma \triangleright_c \varphi \iff \Gamma \triangleright_i \neg\neg\varphi.$$

Proof. Since $\varphi \rightarrow \neg\neg\psi \triangleright_i \neg\neg(\varphi \rightarrow \psi)$ holds, e.g., by [113, lemma 6.2.2], in view of Lemma 5.1(i) $R\rightarrow$ is compatible with $\neg\neg$, and we get the claim as an instance of Corollary 3.7. \square

That $\neg\neg$ is a $\neg\neg$ -translation (Corollary 3.7(b)) is the alternative formulation of Glivenko’s theorem

$$\Gamma \triangleright_c \varphi \implies \neg\neg\Gamma \triangleright_i \neg\neg\varphi.$$

5.2.2. Glivenko’s theorem in general algebra. We compare our approach with some occurrences of Glivenko’s theorem in universal algebra, starting from Birkhoff’s presentation [10, chap. IX, theorem 16]:

*In any pseudo-complemented distributive lattice L , the correspondence $a \mapsto a^{**}$ is a closure operation in L , and a lattice-homomorphism of L onto the complete Boolean algebra of “closed” elements.*

As far as we see, this falls somewhat short of capturing Glivenko’s theorem: Any nucleus $j : S \rightarrow S$ whatsoever on a frame S is a closure operator and induces a frame homomorphism $S \rightarrow S_j$ onto the frame $S_j = \{a \in S : ja = a\}$ in which the join of $C \subseteq S_j$ is $j(\bigvee C)$, where \bigvee denotes the join in S_j , see, e.g., [60, Subsection II.2.2]. As Birkhoff has noticed, moreover, $(a \vee a^*)^{**} = 1$ for every (not necessarily closed) element a , yet we do not see any trace of Glivenko’s theorem. The situation changes with later adaptations such as Esakia’s [28, Section A.2]:

Let H be a Heyting algebra and $Rg(H)$ the set of all its regular elements, i.e.,

$$Rg(H) = \{\neg a : a \in H\}.$$

It is known that $Rg(H)$ is a Boolean algebra and that the map $h : H \rightarrow Rg(H)$, given by $ha = \neg\neg a$ ($a \in H$), is a surjective homomorphism of the Heyting algebra H onto the Boolean algebra $Rg(H)$.

Rather than that the regular elements form a Boolean algebra [8], the crucial issue for Glivenko’s theorem is that double negation is a morphism of Heyting algebras, more precisely that double negation commutes with implication:

$$\neg\neg(\varphi \rightarrow \psi) \approx_i \neg\neg\varphi \rightarrow \neg\neg\psi. \tag{19}$$

We notice that by Remark 2.3(iv) the left-to-right direction of (19) is a general property of nuclei, whereas the right-to-left direction is equivalent to

$$\varphi \rightarrow \neg\neg\psi \triangleright_i \neg\neg(\varphi \rightarrow \psi), \tag{20}$$

which is nothing but the condition that we have exhibited in Lemma 5.1 for the rule $R\rightarrow$ to be compatible with double negation over intuitionistic provability \triangleright_i .

5.2.3. The minimal case. Take propositional minimal logic \triangleright_m as \triangleright . As observed above, stability (8) equals double negation elimination (16); whence the strong extension $\triangleright_m^{g,\perp}$ of minimal logic \triangleright_m is classical logic \triangleright_c .

LEMMA 5.5. *Over any given extension \triangleright_* of minimal logic, axiom (13), viz.*

$$\triangleright_* \neg\neg(\perp \rightarrow \varphi),$$

is equivalent to (20) with \triangleright_ in place of \triangleright_i ; that is,*

$$\varphi \rightarrow \neg\neg\psi \triangleright_* \neg\neg(\varphi \rightarrow \psi). \tag{21}$$

Proof. We can obtain (13) from the instance

$$\perp \rightarrow \neg\neg\varphi \triangleright_* \neg\neg(\perp \rightarrow \varphi),$$

of (21) by an application of Trans with the derivable $\triangleright_* \perp \rightarrow \neg\neg\varphi$.

As (20) is well-known [113], (21) follows from instances of *ex falso sequitur quodlibet* in minimal logic. By using Lj, (21) equally follows from the same instances of (13). In the Appendix 7 we detail a direct proof. \square

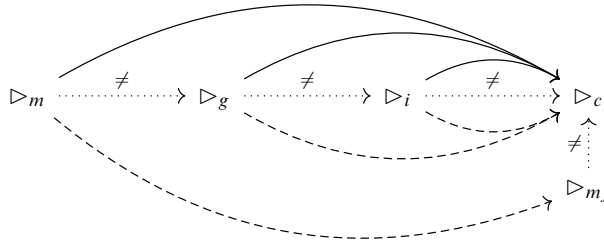


Figure 3. Diagram of the entailment relations involved in the situation of Propositions 5.4 and 5.6. A solid arrow denotes a strong $\neg\neg$ -extension, a dashed arrow denotes a weak $\neg\neg$ -extension, a dotted arrow denotes a generic extension, and a double line denotes a conservative extension.

The following statement [15, 82]⁹ is a generalisation of Proposition 5.4 (Figure 3):

PROPOSITION 5.6 (General Glivenko theorem [15, 82]). *Let \triangleright_* be an inductive extension of \triangleright_m the additional rules of which hold for \triangleright_c . The following are equivalent:*

- (a) $\Gamma \triangleright_c \varphi \iff \Gamma \triangleright_* \neg\neg\varphi$ for all Γ, φ ;
- (b) $\triangleright_g \subseteq \triangleright_*$.

Proof. Since $\varphi \rightarrow \neg\neg\psi \triangleright_i \neg\neg(\varphi \rightarrow \psi)$ holds by 5.5, in view of Lemma 5.1(i) the rule $R\rightarrow$ is compatible with $\neg\neg$, and we get the claim as an instance of Corollary 3.7. □

Since in minimal logic \perp has no special role, we would get a completely analogous statement if we considered the more general continuation nucleus $g_\alpha: \varphi \mapsto (\varphi \rightarrow \alpha) \rightarrow \alpha$ in place of $\neg\neg$ and used the axiom

$$\triangleright ((\alpha \rightarrow \varphi) \rightarrow \alpha) \rightarrow \alpha$$

in place of (13) to define \triangleright_g .

5.2.4. *The predicate case.* Let us go back to considering intuitionistic logic \triangleright_i as \triangleright .

PROPOSITION 5.7 (Gödel’s theorem [46]). *Let \triangleright_* be an inductive extension of \triangleright_i the additional rules of which hold for \triangleright_c . The following are equivalent in predicate logic:*

- (a) $\Gamma \triangleright_c \varphi \iff \Gamma \triangleright_* \neg\neg\varphi$ for all Γ, φ ;
- (b) $\forall x \neg\neg\varphi \triangleright_* \neg\neg\forall x\varphi$ for all φ .

Proof. Similar to Proposition 5.4; but apply Theorem 3.5 instead of Corollary 3.7. Lemma 5.1(ii) can be applied since the Glivenko nucleus is compatible with substitution. □

⁹ This was observed by Odintsov [82] and Cignoli and Torres [15] apparently simultaneously, in 2004. The latter further showed that our conservation criterion $\varphi \rightarrow \neg\neg\psi \triangleright_* \neg\neg(\varphi \rightarrow \psi)$ is equivalent to the double negation of double negation elimination $\triangleright_* \neg\neg(\neg\neg\varphi \rightarrow \varphi)$.

Condition (b) in Proposition 5.7 is called *Double Negation Shift* (DNS) and is known to define a proper intermediate logic $\triangleright_{\text{DNS}}$, that is, $\triangleright_i \subsetneq \triangleright_{\text{DNS}} \subsetneq \triangleright_c$ [31].¹⁰

PROPOSITION 5.8. *Let k, g, J be respectively the Kolmogorov, Gentzen and Kuroda $\neg\neg$ -functions. In predicate logic:*

- (i) $\Gamma \triangleright_c \varphi \iff k\Gamma \triangleright_i k\varphi$.
In other words, k is a $\neg\neg$ -translation, known as the Kolmogorov negative translation [37, 64].
- (ii) $\Gamma \triangleright_c \varphi \iff g\Gamma \triangleright_i g\varphi$.
In other words, g is a $\neg\neg$ -translation, known as the Gentzen negative translation¹¹ [37, 41, 43].
- (iii) $\Gamma \triangleright_c \varphi \iff \neg\neg J\Gamma \triangleright_i \neg\neg J\varphi \iff J\Gamma \triangleright_i \neg\neg J\varphi$.
In other words, $\neg\neg J$ is a $\neg\neg$ -translation, known as the minimal Kuroda negative translation¹² [37, 66, 73, 114].

Proof. Since $\perp \approx_i \neg\neg\perp$, it is easy to see that k, g, J all satisfy (9). Hence by Remark 5.2, k and g satisfy the conditions of Theorem 3.13, which shows (i) and (ii). Similarly, by Remark 5.2, J satisfies the conditions of Corollary 3.14, which shows (iii). \square

Theorem 3.13 gives results similar to those in Proposition 5.8 also for other negative translations, such as the *Gödel negative translation* [47], the original *Kuroda negative translation* [66], the *Krivine negative translation*—which was introduced by Streicher and Reus [108] based on Krivine’s work [65]—and the negative translation by Ferreira and Oliva [37]. We only point out that, unlike the ones presented here, the latter two cannot be generalised—at least not in a direct way—for an arbitrary nucleus j , and hence a counterpart of Remark 5.2 cannot be given.

5.2.5. The infinitary case. The following statement [110]¹³ is the counterpart of Proposition 5.7 for infinitary logic:

PROPOSITION 5.9. *Let \triangleright_* be an inductive extension of \triangleright_i the additional rules of which hold for \triangleright_c . The following are equivalent in infinitary logic:*

- (a) $\Gamma \triangleright_c \varphi \iff \Gamma \triangleright_* \neg\neg\varphi$ for all Γ, φ ;
- (b) $\bigwedge_{i \in \mathbb{N}} \neg\neg\varphi_i \triangleright_* \neg\neg \bigwedge_{i \in \mathbb{N}} \varphi_i$ for all $\{\varphi_i : i \in \mathbb{N}\}$.

Proof. Similar to Proposition 5.7. \square

As above, condition (b) in Proposition 5.9 defines a proper intermediate logic $\triangleright_{\text{DNS}_{\text{inf}}}$, that is, $\triangleright_i \subsetneq \triangleright_{\text{DNS}_{\text{inf}}} \subsetneq \triangleright_c$.

¹⁰ Even though DNS is not valid in intuitionistic logic, it holds constructive value, as noted by Ilik [54].

¹¹ In literature, this is often referred to as the Gödel–Gentzen negative translation. However, the translation introduced by Gödel [47] was in fact a different one, somewhere in between Kolmogorov’s and Gentzen’s [33, 37, 112].

¹² This is a variant introduced by Murthy [73] of the original Kuroda negative translation [66]. It has been studied in the literature for having somewhat nicer properties than Kuroda’s original version [37, 114].

¹³ We are grateful to Matteo Tesi for the advance communication of his result.

5.2.6. *The Peirce nucleus and the Clavius nucleus.* Let again \triangleright_* be an inductive extension of \triangleright_m , and fix a formula α . The *Peirce nucleus* [60, 62, 93, 114] is

$$p_\alpha: \varphi \mapsto (\varphi \rightarrow \alpha) \rightarrow \varphi.$$

As for Lemma 5.3 one readily proves with the generating rules of \triangleright_m that p_α is in fact a nucleus. We are especially interested in the particular case of the *Clavius nucleus*:

$$p_\perp: \varphi \mapsto \neg\varphi \rightarrow \varphi.$$

Over intuitionistic logic, it is easy to see that the Glivenko nucleus is equivalent to the Clavius nucleus, i.e., $\neg\neg\varphi \approx_i \neg\varphi \rightarrow \varphi$ for every φ , and hence Propositions 5.4 and 5.7–5.9 equally hold for the Clavius nucleus in place of the Glivenko nucleus.

If we consider minimal logic \triangleright_m instead, then the Clavius nucleus is no more equivalent to the Glivenko nucleus. As stability (8) for the Clavius nucleus j equals (14), the strong extension \triangleright_m^j of minimal logic \triangleright_m is nothing but the Clavius logic \triangleright_s . We thus get the following counterpart of Glivenko’s theorem:

PROPOSITION 5.10. *In propositional logic,*

$$\Gamma \triangleright_s \varphi \iff \Gamma \triangleright_m \neg\varphi \rightarrow \varphi.$$

Proof. Since $\varphi \rightarrow (\neg\psi \rightarrow \psi) \triangleright_m \neg(\varphi \rightarrow \psi) \rightarrow (\varphi \rightarrow \psi)$, by Lemma 5.1(i) we get the claim as an instance of Corollary 3.7. \square

5.3. *Open and closed nuclei.* Once more let \triangleright_* be an inductive extension of \triangleright_m , and fix a formula α . The *closed nucleus* c_α and the *open nucleus* o_α [93, 114] are defined by¹⁴

$$c_\alpha: \varphi \mapsto \varphi \vee \alpha, \qquad o_\alpha: \varphi \mapsto \alpha \rightarrow \varphi.$$

Still as for Lemma 5.3, with the generating rules of \triangleright_m one easily sees that c_α and o_α are nuclei over \triangleright_* . We notice the following facts about the extensions induced by these nuclei:

- Stability of the open nucleus is equivalent to $\triangleright \alpha$, which can be read as “ α is derivable”; thus the strong extension $\triangleright_*^{o_\alpha}$ is the smallest entailment relation containing \triangleright_* in which “ α holds”. If $\alpha = \perp$, then stability of o_\perp becomes (18), thus the strong extension $\triangleright_m^{o_\perp}$ is nothing but negative logic \triangleright_n .

¹⁴ Why “open” and “closed”? The following is well-known [60]. Let L be a bounded distributive lattice, and let $X = \text{Spec}(L)$ be the *spectrum* of L (i.e., the collection of prime filters of L). Then X is a topological space with basis of opens

$$\text{ext}(a) = \{P \in X : a \in P\} \qquad (a \in L).$$

Consider $a \in L$, write $U = \text{ext}(a)$. Then $X \setminus U = \text{Spec}(L^U)$ and $U = \text{Spec}(L_U)$, where L^U and L_U are L with \leq modified:

$$\begin{aligned} x \leq^U y &\iff x \leq y \vee a, && \text{(closed case),} \\ x \leq_U y &\iff x \wedge a \leq y, && \text{(open case).} \end{aligned}$$

For instance, $a \leq^U 0$ and $1 \leq_U a$. If L is a Heyting algebra, then $x \wedge a \leq y \iff x \leq a \rightarrow y$.

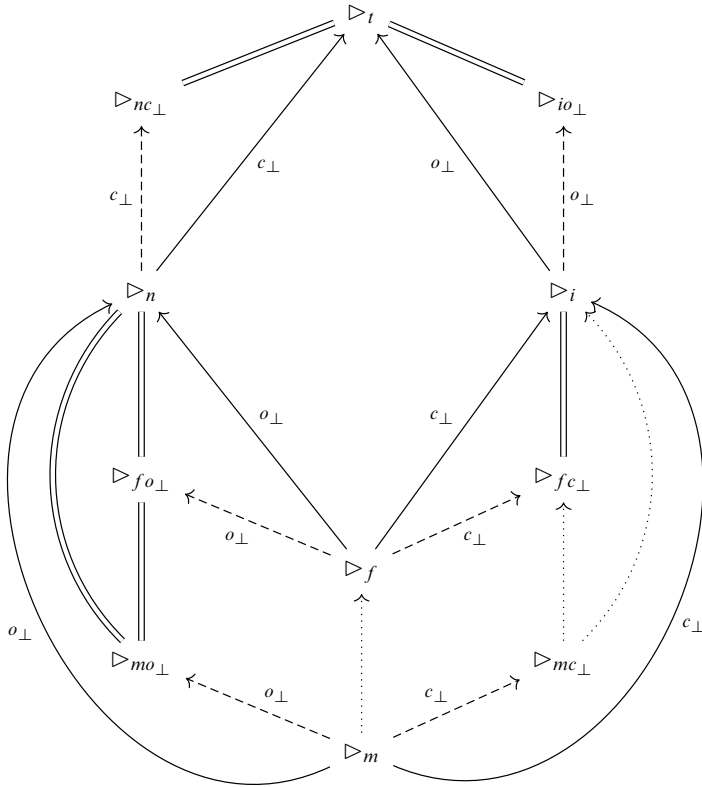


Figure 4. Diagram of the entailment relations induced by open and closed nuclei. A solid arrow denotes a strong extension, a dashed arrow denotes a weak extension, a dotted arrow denotes a generic extension, and a double line denotes a conservative extension.

- Similarly, stability of the closed nucleus is equivalent to $\alpha \triangleright \varphi$ for every propositional formula φ , which can be read as “ α is inconsistent”; thus the strong extension $\triangleright_*^{c\alpha}$ is the smallest entailment relation containing \triangleright_* in which “ α is inconsistent”. If $\alpha = \perp$, then stability of c_\perp becomes (12), thus the strong extension $\triangleright_m^{c\perp}$ is nothing but intuitionistic logic \triangleright_i .
- We also note that in $(\triangleright^{o\alpha})^{c\alpha}$ and $(\triangleright^{c\alpha})^{o\alpha}$ we have both $\triangleright\alpha$ and $\alpha \triangleright \varphi$, which by Trans lead to $\triangleright\varphi$ for every φ : this means that they both equal trivial logic \triangleright_t .

We can sum up the situation via the diagram in Figures 4 and 5. We now ask ourselves whether some of the strong extensions are conservative over the corresponding weak extension.

PROPOSITION 5.11.

(i) Let \triangleright_* be \triangleright_m plus additional axioms. Then:

$$\Gamma \triangleright_*^{o\alpha} \varphi \iff \Gamma \triangleright_* \alpha \rightarrow \varphi \iff \Gamma, \alpha \triangleright_* \varphi,$$

for all Γ, φ . In other words, φ is derivable from Γ when assuming that α is derivable, if and only if $\alpha \rightarrow \varphi$ is derivable from Γ , if and only if φ is derivable from Γ and α .

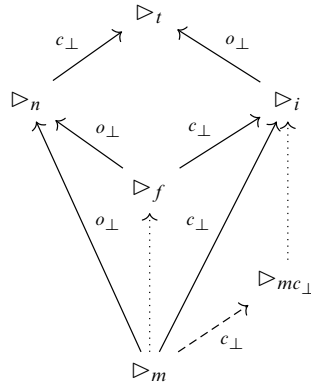


Figure 5. Simplified version of Figure 4, where the weak extensions that coincide with the corresponding strong extension are omitted.

(ii) Let \triangleright_* be \triangleright_m plus additional axioms. Then

$$\Gamma \triangleright_*^{c\alpha} \varphi \iff \Gamma \triangleright_* \varphi \vee \alpha$$

for all Γ, φ if and only if

$$\triangleright_* \alpha \vee (\alpha \rightarrow \varphi). \tag{22}$$

Proof.

- (i) In view of $\varphi \rightarrow (\alpha \rightarrow \psi) \triangleright_m \alpha \rightarrow (\varphi \rightarrow \psi)$, the first equivalence holds by means of Theorem 3.5 and Lemma 5.1. The second equivalence is the *deduction theorem* for \triangleright_* .
- (ii) It can easily be shown that (22) is equivalent to $\varphi \rightarrow (\psi \vee \alpha) \triangleright (\varphi \rightarrow \psi) \vee \alpha$; hence the claim follows from Theorem 3.5 and Lemma 5.1. \square

COROLLARY 5.12.

- (i) Let $\triangleright_* \in \{\triangleright_m, \triangleright_f, \triangleright_i\}$. Then $\triangleright_*^{o\perp} = \triangleright_{*o\perp}$.
- (ii) Let $\triangleright_* \in \{\triangleright_f, \triangleright_n\}$. Then $\triangleright_{*c\perp} = \triangleright_*^{c\perp}$.
- (iii) $\triangleright_{mc\perp} \subsetneq \triangleright_c$.

In conclusion, all strong extensions in Figure 4 are conservative over the corresponding weak extension except for \triangleright_c over $\triangleright_{mc\perp}$.

5.3.1. *A few remarks about the predicate case.* In predicate logic, we get an analogous situation by considering the logic \triangleright_F , defined as the predicate version of the Frobenius logic \triangleright_f plus the *dual Frobenius rule* [62]

$$\forall x(\varphi \vee \perp) \triangleright (\forall x\varphi) \vee \perp.$$

Functions based on the open and closed nuclei but over predicate intuitionistic logic \triangleright_i are known in literature. For instance, having fixed a formula A , Friedman employed the prime c_A -function¹⁵ t_α^C —thus known as *Friedman’s A-translation*—to

¹⁵ See §5.1.

prove Markov’s rule [38]; and Ishihara and Nemoto used the prime o_A -function t_A^O to prove the independence-of-premiss rule [56]. Note that:

$$t_A^O o_A \varphi = t_A^O A \rightarrow t_A^O \varphi, \quad t_A^C c_A \varphi = t_A^C \varphi \vee t_A^C A.$$

This means that t_A^O satisfies (8), and hence is an o_A -translation, if and only if “ $t_A^O A$ holds” in the sense that $\triangleright t_A^O A$; this is the case, e.g., when A is atomic. Similarly, t_A^C satisfies (8), and thus is a c_A -translation, precisely when “ $t_A^C A$ is inconsistent” in the sense that $t_A^C A \triangleright \varphi$ for every φ ; this is the case, e.g., when $A = \perp$. In the latter case, however, the strong c_A -extension is trivial.

5.4. Propositional lax logic. In propositional lax logic [32] the modality \bigcirc is characterised by rules corresponding [32, p. 2, Section 1] to the ones of a nucleus. In our setting, we define \triangleright_L by adding a unary relation symbol \bigcirc to the language of \triangleright_i , and the axiom

$$\varphi \rightarrow \bigcirc \psi \approx \bigcirc \varphi \rightarrow \bigcirc \psi$$

to the inductive definition of \triangleright_i . It is evident that $j = \bigcirc$ is a nucleus over \triangleright_L . This trivially extends to every inductive extension of \triangleright_L since $L\bigcirc$ is a generating rule.

The strong \bigcirc -extension $\triangleright_L^\bigcirc$ of \triangleright_L is \triangleright_i plus

$$\varphi \approx \bigcirc \varphi,$$

i.e., plus an *identity operator* \bigcirc on formulae. We also define \triangleright_{L^*} as \triangleright_L plus

$$\varphi \rightarrow \bigcirc \psi \triangleright \bigcirc(\varphi \rightarrow \psi),$$

i.e., the smallest extension of \triangleright_L in which the nucleus \bigcirc is compatible with $R \rightarrow$ (Theorem 3.5).

Given a formula φ in the language of \triangleright_L , we denote by φ' the formula obtained from φ by removing all occurrences of \bigcirc . Formally, the definition of φ' is given inductively:

$$\begin{aligned} P' &= P, & \top' &= \top, & \perp' &= \perp, \\ (\varphi \wedge \psi)' &= \varphi' \wedge \psi', & (\varphi \vee \psi)' &= \varphi' \vee \psi', & (\varphi \rightarrow \psi)' &= \varphi' \rightarrow \psi', \\ (\forall x \varphi)' &= \forall x \varphi', & (\exists x \varphi)' &= \exists x \varphi', & (\bigcirc \varphi)' &= \varphi'. \end{aligned}$$

We obtain a somewhat more general version of *strong conservativity* [32, theorem 2.4]:

PROPOSITION 5.13. *The following are equivalent:*

- (a) $\Gamma \triangleright_{L^*} \bigcirc \varphi$,
- (b) $\Gamma \triangleright_L^\bigcirc \varphi$,
- (c) $\Gamma' \triangleright_L \varphi'$,
- (d) $\Gamma' \triangleright_i \varphi'$.

In particular, if $\Gamma \triangleright_L \varphi$, then $\Gamma' \triangleright_i \varphi'$.

Proof. Items (a) and (b) are equivalent by Theorem 3.5. The fact that (d) implies (c) holds since $\triangleright_L^\bigcirc$ extends \triangleright_i . The directions from (c) to (b) and from (b) to (d) are proved straightforwardly by structural induction. \square

§6. Future work. Orevkov [84] has established some well-known conservativity results of classical logic over intuitionistic and minimal first-order logics with equality.

In particular, he isolates seven classes of single-succedent sequents—the so-called *Glivenko sequent classes*—defined in terms of the absence of positive or negative occurrences of particular logical symbols (in a first-order language with equality) where classical derivability implies intuitionistic derivability. The same article also shows that these classes are optimal: any class of sequents for which classical derivability implies intuitionistic derivability is contained in one of these seven classes. In recent years simpler proofs of conservativity results for some Glivenko sequent classes have been given [55, 74, 100]. An extremely simple and purely logical proof of the first Glivenko class for coherent theories has been obtained by Negri [77] by means of **G3**-style sequent calculi; this has been extended to geometric theories in [79] and then to cover all other first-order Glivenko sequent classes in [34, 78]. We want to investigate whether a generalisation of these results is possible in our setting.

§7. Appendix Proof details. We give here a direct proof of Lemma 5.5.

LEMMA 7.1. *Over any given extension \triangleright_* of minimal logic, the following are equivalent:*

- (a) $\triangleright_* \neg(\perp \rightarrow \varphi)$,
 (b) $\varphi \rightarrow \neg\neg\psi \triangleright_* \neg(\varphi \rightarrow \psi)$.

Proof. As for (b) \implies (a):

$$\frac{\frac{\frac{\text{Refl}}{\perp, \neg\varphi \triangleright_* \perp}}{\perp \triangleright_* \neg\neg\varphi} \text{R}\rightarrow}{\triangleright_* \perp \rightarrow \neg\neg\varphi} \text{R}\rightarrow \quad \frac{\text{assumption}}{\perp \rightarrow \neg\neg\varphi \triangleright_* \neg(\perp \rightarrow \varphi)} \text{Trans}}{\triangleright_* \neg(\perp \rightarrow \varphi)} \text{Trans}$$

As for (a) \implies (b):

$$\frac{\frac{\frac{\text{axiom}}{\varphi \rightarrow \neg\neg\psi, \varphi \triangleright_* \neg\neg\psi} \text{L}\neg\neg}{\varphi \rightarrow \neg\neg\psi, \neg\neg\varphi \triangleright_* \neg\neg\psi} \text{R}\rightarrow}{\varphi \rightarrow \neg\neg\psi \triangleright_* \neg\neg\varphi \rightarrow \neg\neg\psi} \text{R}\rightarrow \quad \frac{\text{(see below)} \quad \frac{\frac{\text{Refl}}{\psi, \neg(\varphi \rightarrow \psi) \triangleright_* \perp} \text{L}\neg\neg}{\psi, \neg(\varphi \rightarrow \psi) \triangleright_* \perp} \text{L}\rightarrow}{\psi, \neg(\varphi \rightarrow \psi) \triangleright_* \perp} \text{L}\rightarrow}{\neg(\varphi \rightarrow \psi) \triangleright_* \neg\neg\varphi} \text{L}\rightarrow}{\varphi \rightarrow \neg\neg\psi \triangleright_* \neg(\varphi \rightarrow \psi)} \text{Trans}$$

where $\neg(\varphi \rightarrow \psi) \triangleright_* \neg\neg\varphi$ is derived as follows:

$$\frac{\frac{\frac{\text{Refl}}{\perp \rightarrow \psi, \perp, \neg\varphi \triangleright_* \perp} \text{Refl}}{\perp \rightarrow \psi, \neg(\varphi \rightarrow \psi), \neg\varphi \triangleright_* \perp} \text{R}\rightarrow \quad \frac{\frac{\text{axiom}}{\psi, \neg\varphi, \varphi \triangleright_* \psi} \text{Refl}}{\psi, \neg\varphi, \varphi \triangleright_* \psi} \text{L}\rightarrow}{\perp \rightarrow \psi, \neg\varphi, \varphi \triangleright_* \psi} \text{R}\rightarrow}{\perp \rightarrow \psi, \neg(\varphi \rightarrow \psi), \neg\varphi \triangleright_* \perp} \text{R}\rightarrow}{\perp \rightarrow \psi, \neg(\varphi \rightarrow \psi) \triangleright_* \neg\neg\varphi} \text{L}\rightarrow}{\frac{\text{(a)} \quad \frac{\frac{\text{Refl}}{\perp \rightarrow \psi, \neg(\varphi \rightarrow \psi)} \text{L}\rightarrow}{\perp \rightarrow \psi, \neg(\varphi \rightarrow \psi) \triangleright_* \neg\neg\varphi} \text{L}\rightarrow}{\triangleright_* \neg\neg(\perp \rightarrow \psi)} \text{Trans}}{\neg(\varphi \rightarrow \psi) \triangleright_* \neg\neg\varphi} \text{Trans} \quad \square$$

Acknowledgements. The authors wish to thank Benno van den Berg, Ingo Blechschmidt, Hajime Ishihara, Sara Negri, Paulo Oliva, Mario Piazza, Matteo Tesi

and Tarmo Uustalu for their interest and suggestions. Last but not least, the authors are grateful to Daniel Misselbeck-Wessel whose original idea set the basis for this work.

Funding. The present study was carried out within the projects “A New Dawn of Intuitionism: Mathematical and Philosophical Advances” (ID 60842) funded by the John Templeton Foundation, and “Reducing complexity in algebra, logic, combinatorics - REDCOM” belonging to the programme “Ricerca Scientifica di Eccellenza 2018” of the Fondazione Cariverona. Both authors are members of the “Gruppo Nazionale per le Strutture Algebriche, Geometriche e le loro Applicazioni” (GNSAGA) of the Istituto Nazionale di Alta Matematica (INdAM).

BIBLIOGRAPHY

- [1] Aczel, P. (1978). The type theoretic interpretation of constructive set theory. In *Logic Colloquium '77 (Proc. Conf., Wrocław, 1977)*, Studies in Logic and the Foundations of Mathematics, 96. Amsterdam: North-Holland, pp. 55–66.
- [2] ———. (1982). The type theoretic interpretation of constructive set theory: Choice principles. In *The L. E. J. Brouwer Centenary Symposium (Noordwijkerhout, 1981)*, Studies in Logic and the Foundations of Mathematics, 110. Amsterdam: North-Holland, pp. 1–40.
- [3] ———. (1986). The type theoretic interpretation of constructive set theory: inductive definitions. In *Logic, Methodology and Philosophy of Science, VII (Salzburg, 1983)*, Studies in Logic and the Foundations of Mathematics, 114. Amsterdam: North-Holland, pp. 17–49.
- [4] ———. (2001). The Russell–Prawitz modality. *Mathematical Structures in Computer Science*, **11**(4), 541–554.
- [5] ———. (2006). Aspects of general topology in constructive set theory. *Annals of Pure and Applied Logic*, **137**(1–3), 3–29.
- [6] Aczel, P., & Rathjen, M. (2000). Notes on constructive set theory. Technical report, Institut Mittag-Leffler (The Royal Swedish Academy of Sciences), Report No. 40.
- [7] ———. (2010). *Constructive Set Theory*. Book draft. <https://ncatlab.org/nlab/files/AczelRathjenCST.pdf>.
- [8] Bezhanišvili, G., & Holliday, W. H. (2019). A semantic hierarchy for intuitionistic logic. *Indagationes Mathematicae*, **30**(3), 403–469.
- [9] Béziau, J.-Y. (2006). Les axiomes de Tarski. In Pouivet, R., & Resbuschi, M., editors. *La Philosophie en Pologne 1919–1939*. Paris: Librairie Philosophique J. VRIN, pp. 135–149.
- [10] Birkhoff, G. (1948). *Lattice Theory*, American Mathematical Society Colloquium Publications, 25 (revised edition). New York: American Mathematical Society.
- [11] Bishop, E., & Bridges, D. (1985). *Constructive Analysis*. Berlin: Springer.
- [12] ———. (1987). Constructive analysis. *Journal of Symbolic Logic*, **52**(4), 1047–1048.
- [13] Bishop, E. A. (1967). *Foundations of Constructive Analysis*. New York: McGraw-Hill.
- [14] Cederquist, J., & Coquand, T. (2000). Entailment relations and distributive lattices. In Buss, S. R., Hájek, P., & Pudlák, P., editors. *Logic Colloquium '98. Proceedings of the Annual European Summer Meeting of the Association for Symbolic*

Logic, Prague, Czech Republic, August 9–15, 1998. Lecture Notes in Logic, 13. Natick: A. K. Peters, pp. 127–139.

[15] Cignoli, R., & Torrens, A. (2004). Glivenko like theorems in natural expansions of BCK-logic. *Mathematical Logic Quarterly*, **50**, 111–125.

[16] Cintula, P., & Carles, N. (2013). The proof by cases property and its variants in structural consequence relations. *Studia Logica*, **101**(4), 713–747.

[17] Ciraulo, F., Maietti, M. E., & Sambin, G. (2013). Convergence in formal topology: A unifying notion. *Journal of Logic and Analysis*, **5**(2), 1–45.

[18] Ciraulo, F., & Sambin, G. (2008). Finitary formal topologies and Stone’s representation theorem. *Theoretical Computer Science*, **405**(1–2), 11–23.

[19] Coquand, T. (2000). *A Direct Proof of the Localic Hahn–Banach Theorem*. Available from: <https://ncatlab.org/nlab/files/AczelRathjenCST.pdf>.

[20] ———. (2000). *Lewis Carroll, Gentzen and Entailment Relations*. Available from: <https://ncatlab.org/nlab/files/AczelRathjenCST.pdf>.

[21] ———. (2005). About Stone’s notion of spectrum. *Journal of Pure and Applied Algebra*, **197**(1–3), 141–158.

[22] ———. (2009). Space of valuations. *Annals of Pure and Applied Logic*, **157**, 97–109.

[23] Coquand, T., & Lombardi, H. (2002). Hidden constructions in abstract algebra: Krull dimension of distributive lattices and commutative rings. In Fontana, M., Kabbaj, S.-E., & Wiegand, S., editors. *Commutative Ring Theory and Applications*, Lecture Notes in Pure and Applied Mathematics, 231. Reading, MA: Addison-Wesley, pp. 477–499.

[24] Coquand, T., Lombardi, H., & Neuwirth, S. (2019). Lattice-ordered groups generated by an ordered group and regular systems of ideals. *The Rocky Mountain Journal of Mathematics*, **49**(5), 1449–1489.

[25] Coquand, T., & Neuwirth, S. (2017). An introduction to Lorenzen’s “algebraic and logistic investigations on free lattices” (1951). Preprint.

[26] Coquand, T., & Persson, H. (2001). Valuations and Dedekind’s Prague theorem. *Journal of Pure and Applied Algebra*, **155**(2–3), 121–129.

[27] Coquand, T., & Zhang, G.-Q. (2000). Sequents, frames, and completeness. In Clote, P. G., & Schwichtenberg, H., editors. *Computer Science Logic (Fischbachau, 2000)*, Lecture Notes in Computer Science, 1862. Berlin: Springer, pp. 277–291.

[28] Esakia, L. L. (2019). *Heyting Algebras: Duality Theory*. Cham: Springer.

[29] Escardó, M., & Oliva, P. (2010). The Peirce translation and the double negation shift. In Ferreira, F., Löwe, B., Mayordomo, E., & Gomes, L. M., editors. *Programs, Proofs, Processes*. Berlin, Heidelberg: Springer, pp. 151–161.

[30] ———. (2012). The Peirce translation. *Annals of Pure and Applied Logic*, **163**, 681–692.

[31] Espindola, C. (2013). A short proof of Glivenko theorems for intermediate predicate logics. *Archive for Mathematical Logic*, **52**(7–8), 823–826.

[32] Fairtlough, M., & Mendler, M. (1997). Propositional lax logic. *Information and Computation*, **137**(1), 1–33.

[33] Fellin, G. (2022). *Constructivisation through Induction and Conservation*. Ph.D. Thesis, University of Helsinki & University of Trento, Unigrafia, Helsinki.

[34] Fellin, G., Negri, S., & Orlandelli, E. (2023). Glivenko sequent classes and constructive cut elimination in geometric logics. *Archive for Mathematical Logic*, **62**, 657–688.

- [35] Fellin, G., & Schuster, P. (2021). A general Glivenko–Gödel theorem for nuclei. In Sokolova, A., editor. *Proceedings of the 37th Conference on the Mathematical Foundations of Programming Semantics, MFPS 2021, Salzburg, Austria, August 29–September 3, 2021*, Electronic Notes in Theoretical Computer Science. Amsterdam: Elsevier, pp. 51–66.
- [36] Fellin, G., Schuster, P., & Wessel, D. (2022). The Jacobson radical of a propositional theory. *Bulletin of Symbolic Logic*, **28**(2), 163–181.
- [37] Ferreira, G., & Oliva, P. (2013). On the relation between various negative translations. In Berger, U., Diener, H., Schuster, P., & Seisenberger, M., editors. *Logic, Construction, Computation*. Frankfurt: De Gruyter, pp. 227–258.
- [38] Friedman, H. (1978). Classically and intuitionistically provably recursive functions. In Müller, G. H., & Scott, D. S., editors. *Higher Set Theory (Proc. Conf., Math. Forschungsinst., Oberwolfach, 1977)*. Lecture Notes in Mathematics, 669. Berlin: Springer, pp. 21–27.
- [39] Gabbay, D. M. (1972). Applications of trees to intermediate logics. *Journal of Symbolic Logic*, **37**(1), 135–138.
- [40] Galatos, N., & Ono, H. (2006). Glivenko theorems for substructural logics over FL. *The Journal of Symbolic Logic*, **71**(4), 1353–1384.
- [41] Gentzen, G. (1936). Die Widerspruchsfreiheit der reinen Zahlentheorie. *Mathematische Annalen*, **112**(1), 493–565.
- [42] ———. (1938). Neue Fassung des Widerspruchsfreiheitsbeweises für die reine Zahlentheorie. *Forschungen zur Logik und zur Grundlegung der exakten Wissenschaften*, **4**, 19–44.
- [43] ———. (1974). Über das Verhältnis zwischen intuitionistischer und klassischer Arithmetik. *Archiv für mathematische Logik und Grundlagenforschung*, **16**, 119–132.
- [44] Glivenko, V. (1928). Sur la logique de M. Brouwer. *Académie Royale de Belgique, Bulletins de la classe des sciences (5)*, **14**, 225–228.
- [45] ———. (1929). Sur quelques points de la Logique de M. Brouwer. *Académie Royale de Belgique, Bulletins de la classe des sciences (5)*, **15**, 183–188.
- [46] Gödel, K. (1933). Zur intuitionistischen Arithmetik und Zahlentheorie. *Ergebnisse eines math. Kolloquiums*, **4**, 39–40.
- [47] ———. (1986). Zur intuitionistischen Arithmetik und Zahlentheorie (gekuerzter Nachdruck). In Berka, K., & Kreiser, L., editors. *Logik-Texte: Kommentierte Auswahl Zur Geschichte der Modernen Logik (Vierte Auflage)*. Berlin: Akademie-Verlag, pp. 201–202.
- [48] Griffin, T. (1990). A formulae-as-types notion of control. In Allen, F. E., editor. *Conference Record of the Seventeenth Annual ACM Symposium on Principles of Programming Languages, San Francisco, California, USA, January 1990*. New York: ACM Press, pp. 47–58. <https://doi.org/10.1145/96709.96714>.
- [49] Guerrieri, G., & Naibo, A. (2019). Postponement of raa and Glivenko’s theorem, revisited. *Studia Logica*, **107**(1), 109–144.
- [50] Haykazyan, L. (2020). More on a curious nucleus. *Journal of Pure and Applied Algebra*, **224**, 860–868.
- [51] Hertz, P. (1922). Über Axiomensysteme für beliebige Satzsysteme. I. Teil. Sätze ersten grades. *Mathematische Annalen*, **87**(3), 246–269.
- [52] ———. (1923). Über Axiomensysteme für beliebige Satzsysteme. II. Teil. Sätze höheren grades. *Mathematische Annalen*, **89**(1), 76–102.

- [53] ———. (1929). Über Axiomensysteme für beliebige Satzsysteme. *Mathematische Annalen*, **101**(1), 457–514.
- [54] Ilik, D. (2013). Double-negation shift as a constructive principle. *CoRR*, Preprint, [abs/1301.5089](https://arxiv.org/abs/1301.5089).
- [55] Ishihara, H. (2000). A note on the Gödel–Gentzen translation. *Mathematical Logic Quarterly*, **46**, 135–137.
- [56] Ishihara, H., & Nemoto, T. (2016). A note on the independence of premiss rule. *Mathematical Logic Quarterly*, **62**(1–2), 72–76.
- [57] Ishihara, H., & Schwichtenberg, H. (2016). Embedding classical in minimal implicational logic. *Mathematical Logic Quarterly*, **62**(1–2), 94–101.
- [58] Jibladze, M., & Johnstone, P. T. (1991). The frame of fibrewise closed nuclei. *Cahiers de Topologie et Geometrie Differentielle Categoriqes*, **32**, 99–112.
- [59] Johansson, I. (1937). Der Minimalkalkül, ein reduzierter intuitionistischer Formalismus. *Compositio Mathematica*, **4**, 119–136.
- [60] Johnstone, P. T. (1982). *Stone Spaces*, Cambridge Studies in Advanced Mathematics, 3. Cambridge: Cambridge University Press.
- [61] ———. (1991). The art of pointless thinking: a student’s guide to the category of locales. In Herrlich, H., & Porst, H.-E., editors. *Category Theory at Work*. Berlin: Heldermann Verlag, pp. 85–107.
- [62] ———. (2002). *Sketches of an Elephant: A Topos Theory Compendium*, Vol. 1 & 2. New York: Oxford University Press.
- [63] Kohlenbach, U. (2008). *Applied Proof Theory: Proof Interpretations and Their Use in Mathematics*, Springer Monographs in Mathematics. Berlin: Springer-Verlag.
- [64] Kolmogorov, A. N. (1925). On the principle of the excluded middle (Russian). *Matematicheskii Sbornik*, **32**, 646–667.
- [65] Krivine, J.-L. (1990). Opérateurs de mise en mémoire et traduction de Gödel. *Archive for Mathematical Logic*, **30**(4), 241–267.
- [66] Kuroda, S. (1951). Intuitionistische Untersuchungen der formalistischen Logik. *Nagoya Mathematical Journal*, **3**, 35–47.
- [67] Legris, J. (2012). Paul Hertz and the origins of structural reasoning. In Béziau, J.-Y., editor. *Universal Logic: An Anthology. From Paul Hertz to Dov Gabbay*, Studies in Universal Logic. Basel: Birkhäuser, pp. 3–10.
- [68] Litak, T., Polzer, M., & Rabenstein, U. (2017). Negative translations and normal modality. In *2nd International Conference on Formal Structures for Computation and Deduction*, LIPIcs, Leibniz International Proceedings in Informatics, 84. Article no. 27, 18. Wadern: Schloss Dagstuhl. Leibniz-Zent. Inform.
- [69] Lombardi, H., & Quitté, C. (2015). *Commutative Algebra: Constructive Methods. Finite Projective Modules*, Algebra and Applications, 20. Dordrecht: Springer.
- [70] Macnab, D. S. (1981). Modal operators on heyting algebras. *Algebra Universalis*, **12**, 5–29.
- [71] Mines, R., Richman, F., & Ruitenburg, W. (1988). *A Course in Constructive Algebra*. New York: Springer.
- [72] Miquel, A. (2020). Implicative algebras: A new foundation for realizability and forcing. *Mathematical Structures in Computer Science*, **30**(5), 458–510.
- [73] Murthy, C. R. (1990). *Extracting Constructive Content from Classical Proofs*. Ph.D. Thesis, Cornell University.

[74] Nadathur, G. (2000). Correspondences between classical, intuitionistic and uniform provability. *Theoretical Computer Science*, **232**, 273–298.

[75] Negri, S. (1996). Stone bases alias the constructive content of stone representation. In Ursini, A., & Aglianò, P., editors. *Logic and Algebra. Proceedings of the International Conference Dedicated to the Memory of Roberto Magari, April 26–30, 1994, Pontignano, Italy*, Lecture Notes in Pure and Applied Mathematics, 180. New York: Marcel Dekker, pp. 617–636.

[76] ———. (2002). Continuous domains as formal spaces. *Mathematical Structures in Computer Science*, **12**(1), 19–52.

[77] ———. (2003). Contraction-free sequent calculi for geometric theories with an application to Barr's theorem. *Archive for Mathematical Logic*, **42**(4), 389–401.

[78] ———. (2016). Glivenko sequent classes in the light of structural proof theory. *Archive for Mathematical Logic*, **55**(3–4), 461–473.

[79] ———. (2021). Geometric rules in infinitary logic. In Arieli, O., & Zamansky, A., editors. *Arnon Avron on Semantics and Proof Theory of Non-Classical Logics*. Cham: Springer, pp. 265–293.

[80] Negri, S., von Plato, J., & Coquand, T. (2004). Proof-theoretical analysis of order relations. *Archive for Mathematical Logic*, **43**, 297–309.

[81] Odintsov, S. (2008). *Constructive Negations and Paraconsistency*. Dordrecht: Springer.

[82] Odintsov, S. P. (2004). Negative equivalence of extensions of minimal logic. *Studia Logica*, **78**(3), 417–442.

[83] Ono, H. (2009). Glivenko theorems revisited. *Annals of Pure and Applied Logic*, **161**(2), 246–250.

[84] Orevkov, V. P. (1968). Glivenko's sequence classes. In Orevkov, V. P., editor. *Logical and Logico-mathematical Calculi. Part I*. Providence: American Mathematical Society, pp. 131–154.

[85] Pereira, L. C., & Haeusler, E. H. (2015). On constructive fragments of classical logic. In Wansing, H., editor. *Dag Prawitz on Proofs and Meaning*, Outstanding Contributions to Logic, 7. Cham: Springer, pp. 281–292.

[86] Picado, J., & Pultr, A. (2012). *Frames and Locales. Topology without Points*, Frontiers in Mathematics. Basel: Birkhäuser.

[87] Restall, G. (2000). *An Introduction to Substructural Logics*. London: Routledge.

[88] Rinaldi, D. (2014). Formal Methods in the Theories of Rings and Domains. Doctoral dissertation, Universität München.

[89] Rinaldi, D., Schuster, P., & Wessel, D. (2017). Eliminating disjunctions by disjunction elimination. *Bulletin of Symbolic Logic*, **23**(2), 181–200.

[90] ———. (2018). Eliminating disjunctions by disjunction elimination. *Indagationes Mathematicae (N.S.)*, **29**(1), 226–259. Communicated first in *Bulletin of Symbolic Logic*, 23(2017), 181–200.

[91] Rinaldi, D., & Wessel, D. (2018). Extension by conservation. Sikorski's theorem. *Logical Methods in Computer Science*, **14**(4–8), 1–17.

[92] ———. (2019). Cut elimination for entailment relations. *Archive for Mathematical Logic*, **58**(5–6), 605–625.

[93] Rosenthal, K. I. (1990). *Quantales and Their Applications*, Pitman Research Notes in Mathematics, 234. Essex: Longman Scientific & Technical.

[94] Rump, W. (2009). A general Glivenko theorem. *Algebra Universalis*, **61**, 455–473.

- [95] Sambin, G. (1987). Intuitionistic formal spaces—A first communication. In Skordev, D., editor. *Mathematical Logic and its Applications, Proc. Adv. Internat. Summer School Conf., Druzhba, Bulgaria, 1986*. New York: Plenum, pp. 187–204.
- [96] ———. (2003). Some points in formal topology. *Theoretical Computer Science*, **305**(1–3), 347–408.
- [97] ———. (forthcoming). *The Basic Picture. Structures for Constructive Topology*, Oxford Logic Guides. Oxford: Clarendon Press. Available from: <https://www.math.unipd.it/~sambin/txt/BPbookContentsOct2006.pdf>.
- [98] Schlagbauer, K., Schuster, P., & Wessel, D. (2019). Der Satz von Hahn–Banach per Disjunktionselimination. *Confluentes Mathematici*, **11**(1), 79–93.
- [99] Schuster, P., & Wessel, D. (2020). Resolving finite indeterminacy: A definitive constructive universal prime ideal theorem. In *Proceedings of the 35th Annual ACM/IEEE Symposium on Logic in Computer Science, LICS '20*. New York: Association for Computing Machinery, pp. 820–830.
- [100] Schwichtenberg, H., & Senjak, C. (2013). Minimal from classical proofs. *Annals of Pure and Applied Logic*, **164**(6), 740–748.
- [101] Schwichtenberg, H., & Wainer, S. S. (2012). *Proofs and Computations, Perspectives in Logic*. Cambridge: Cambridge University Press.
- [102] Scott, D. (1971). On engendering an illusion of understanding. *Journal of Philosophy*, **68**, 787–807.
- [103] ———. (1974). Completeness and axiomatizability in many-valued logic. In Henkin, L., Addison, J., Chang, C. C., Craig, W., Scott, D., & Vaught, R., editors. *Proceedings of the Tarski Symposium (Proc. Sympos. Pure Math., Vol. XXV, University of California, Berkeley, CA, 1971)*. Providence: American Mathematical Society, pp. 411–435.
- [104] Scott, D. S. (1973). Background to formalization. In Leblanc, H., editor. *Truth, Syntax and Modality (Proc. Conf. Alternative Semantics, Temple Univ., Philadelphia, PA, 1970)*, Studies in Logic and the Foundations of Mathematics, 68. Amsterdam: North-Holland, pp. 244–273.
- [105] Simmons, H. (1978). A framework for topology. In Macintyre, A., Pacholski, L., & Paris, J., editors. *Logic Colloquium '77*, Studies in Logic and the Foundations of Mathematics, 96. Amsterdam: North-Holland, pp. 239–251.
- [106] ———. (2010). A curious nucleus. *Journal of Pure and Applied Algebra*, **214**, 2063–2073.
- [107] Spector, C. (1962). Provably recursive functionals of analysis: A consistency proof of analysis by an extension of principles formulated in current intuitionistic mathematics. In *Recursive Function Theory: Proceedings of Symposia in Pure Mathematics*, Vol. 5. Providence: American Mathematical Society, pp. 1–27.
- [108] Streicher, T., & Reus, B. (1998). Classical logic, continuation semantics and abstract machines. *Journal of Functional Programming*, **8**(6), 543–572.
- [109] Tarski, A. (1930). Fundamentale Begriffe der Methodologie der deduktiven Wissenschaften. I. *Monatshefte für Mathematik*, **37**, 361–404.
- [110] Tesi, M. (2020). *Gödel–Gentzen*. Pisa: Scuola Normale Superiore di Pisa.
- [111] Torrens, A. (2008). An approach to Glivenko's theorem in algebraizable logics. *Studia Logica*, **88**, 349–383.
- [112] Troelstra, A. S., & van Dalen, D. (1988). *Constructivism in Mathematics: An Introduction*, Studies in Logic and the Foundations of Mathematics, II. Amsterdam: North-Holland.

- [113] van Dalen, D. (2013). *Logic and Structure* (fifth edition). London: Springer.
- [114] van den Berg, B. (2019). A Kuroda-style j -translation. *Archive for Mathematical Logic*, **58**(5), 627–634.
- [115] ———. (2022). Implicative algebras for modified realizability. Conference presentation abstract, *Continuity, Computability, Constructivity. From Logic to Algorithms*, Padua, Italy: Università degli Studi di Padova. Available from: https://events.math.unipd.it/ccc2022/submissions/CCC2022_paper_2336.pdf.
- [116] Wang, S.-m., & Cintula, P. (2008). Logics with disjunction and proof by cases. *Archive for Mathematical Logic*, **47**(5), 435–446.
- [117] Wessel, D. (2019). Ordering groups constructively. *Communications in Algebra*, **47**(12), 4853–4873.
- [118] ———. (2019). Point-free spectra of linear spreads. In Centrone, S., Negri, S., Sarikaya, D., & Schuster, P., editors. *Mathesis Universalis, Computability and Proof*. Synthese Library. Cham: Springer, pp. 353–374.
- [119] Zdanowski, K. (2009). On second order intuitionistic propositional logic without a universal quantifier. *Journal of Symbolic Logic*, **74**, 157–167.

DIPARTIMENTO DI INFORMATICA
UNIVERSITÀ DEGLI STUDI DI VERONA
STRADA LE GRAZIE 15
37134 VERONA
ITALY

E-mail: giulio.fellin@univr.it, peter.schuster@univr.it

DIPARTIMENTO DI INGEGNERIA DELL'INFORMAZIONE
UNIVERSITÀ DEGLI STUDI DI BRESCIA
VIA BRANZE 38
25123 BRESCIA
ITALY

E-mail: giulio.fellin@unibs.it