

SUFFICIENT CONDITIONS FOR MATCHINGS

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1. Introduction

A graph G is said to possess a *perfect matching* if there is a subgraph of G consisting of disjoint edges which together cover all the vertices of G . Clearly G must then have an even number of vertices. A necessary and sufficient condition for G to possess a perfect matching was obtained by Tutte (3). If S is any set of vertices of G , let $p(S)$ denote the number of components of the graph $G - S$ with an odd number of vertices. Then the condition

$$\text{for all } S, p(S) \leq |S|$$

is both necessary and sufficient for the existence of a perfect matching. A simple proof of this result is given in (1).

We consider certain conditions which are sufficient although not necessary. Roughly speaking, G will have a perfect matching if there are enough edges. For example, if $|V(G)| = n$, n even, where $V(G)$ denotes the set of vertices of G , and if each vertex is of degree $\geq \frac{1}{2}n$, i.e. if each vertex has at least $\frac{1}{2}n$ edges incident with it, then it is almost trivial (see § 3) to show that G has a perfect matching. Instead of looking at each vertex separately, we can put a condition on the vertices collectively. If X denotes any subset of $V(G)$, let

$$\Gamma(X) = \{y \in V(G) : y \text{ is joined by an edge to at least one vertex in } X\}.$$

Following Woodall (4), we define

$$\text{melt}(G) = \max \{c : \forall X \subset V(G), |\Gamma(X)| \geq \min(c|X|, |V(G)|)\}.$$

Thus $\text{melt}(G)$ is the largest number c such that any k vertices of G are collectively adjacent to at least $\min(ck, n)$ vertices. We have already (1) shown that, if n is even,

$$\text{melt}(G) \geq \frac{4}{3} \Rightarrow G \text{ has a perfect matching.} \quad (1)$$

We note that this condition implies that each vertex is of degree $\geq \frac{1}{4}n$. Indeed, we have in general

Lemma. *If $\text{melt}(G) \geq c > 1$, then each vertex of G has degree $\geq \frac{c-1}{c}n$ where $n = |V(G)|$.*

Proof. Suppose there is a vertex v of degree $\leq \frac{c-1}{c}n$. Then there are

$\geq \frac{n}{c}$ vertices none of which is joined by an edge to v . But these vertices must be joined to at least $c \cdot \frac{n}{c} = n$ vertices, a contradiction.

In the next section, we combine the two types of condition above to prove

Theorem 1. *Let G have n vertices, n even. Let c be any fixed number, $\frac{1}{4} \leq c \leq \frac{1}{2}$, and suppose that*

- (i) *each vertex is of degree $\geq cn$,*
- (ii) *$\text{melt}(G) \geq \frac{3-4c}{2-2c}$.*

Then G possesses a perfect matching.

Note 1. $c = \frac{1}{2}$ gives the trivial result mentioned above, and $c = \frac{1}{4}$ gives the result (1).

Note 2. The theorem is also true for other values of c , but if $c > \frac{1}{2}$ condition (i) by itself is sufficient, whereas if $c < \frac{1}{4}$ then condition (ii) by itself suffices.

Note 3. Condition (ii) implies, by the lemma, that each vertex has degree $\geq \frac{1-2c}{3-4c}n$, but this is less than cn if $c > \frac{1}{4}$.

Note 4. The result is best possible. If A, B are graphs let $A+B$ denote the graph obtained by joining every vertex of A to every vertex of B . Take $A = aK_3 \cup bK_1$ and $B = (a+b-2)K_1$ where K_n denotes the complete graph on n vertices.

Following the suggestion of the referee, who is to be thanked for his careful consideration of the original version of this paper, we shall deduce Theorem 1 from the following stronger theorem which is proved along the same lines but more simply.

Theorem 2. *Let G have n vertices, n even, and suppose that*

$$|\Gamma(X)| \geq \min\left(2|X| - \frac{n}{2}, n\right)$$

for all sets X of vertices of G . Then either G has a perfect matching or there exist subsets X, Y of $V(G)$, $X \not\subseteq Y$, such that

$$|X| = \frac{1}{4}(3n-6), |Y| = \frac{1}{4}(3n-2), |\Gamma(X)| = 2|X| - \frac{n}{2}, |\Gamma(Y)| = 2|Y| - \frac{n}{2}.$$

An example of a graph in which the second possibility occurs is $G = 3K_3 + K_1$. Theorem 2 is proved in the next section, but we now show that

Theorem 2 implies Theorem 1. We assume Theorem 2 and the hypotheses of Theorem 1. Let W be any set of vertices of G . If $|W| > (1 - c)n$, then, since the degree of each vertex of G is $\geq cn$, we cannot have a vertex of G joined to no vertex of W . Thus $|\Gamma(W)| = V(G)$. So suppose $|W| \leq (1 - c)n$. Then

$$|\Gamma(W)| \geq \frac{3 - 4c}{2 - 2c} |W| \geq 2 |W| - \frac{n}{2}. \tag{2}$$

It follows from Theorem 2 that G possesses a perfect matching unless there exist two sets X, Y as in Theorem 2. Then, by (2),

$$\frac{3 - 4c}{2 - 2c} |W| = 2 |W| - \frac{n}{2}$$

for $W = X$ and for $W = Y$, giving $|X| = |Y|$, a contradiction.

Theorem 2 is proved in the next section. In the remainder of this paper we shall generalize in one theorem both Theorem 1 and a result of Woodall (4) concerned with the maximum number of disjoint edges in a graph with no perfect matching. Woodall's argument was based on that of (1), and now we in turn extend his result.

2. Proof of Theorem 2

We suppose there is no perfect matching of G . Then by Tutte's theorem there is a set S of vertices of G for which $p(S) > |S|$. Using the fact that $p(S) \equiv |S| \pmod{2}$, we must then have

$$p(S) \geq |S| + 2.$$

Case 1. Suppose that $|S| \geq \frac{1}{4}(n - 6)$. Let m denote the number of 1-components of $G - S$ (i.e. the number of components with just one vertex). Then

$$\begin{aligned} n &\geq |S| + m + 3(p(S) - m) \\ &\geq 4 |S| + 6 - 2m, \end{aligned} \tag{3}$$

whence

$$n - m \leq \frac{3}{2}n - 2 |S| - 3. \tag{4}$$

But, if $m > 0$,

$$\begin{aligned} n - m &\geq |\Gamma(G - S)| \geq 2 |G - S| - \frac{n}{2} \\ &= 2n - 2 |S| - \frac{n}{2}, \end{aligned}$$

whence

$$n - m \geq \frac{3}{2}n - 2 |S|. \tag{5}$$

Since (4) and (5) contradict one another, we must have $m = 0$. Thus, from (3),

$$n \geq 4 |S| + 6$$

i.e.

$$|S| \leq \frac{1}{4}(n-6),$$

whence

$$|S| = \frac{1}{4}(n-6).$$

Equality here implies that each component of $G - S$ must have exactly 3 vertices. If we let X denote the set of vertices in all but one of these components we then have $|X| = \frac{1}{4}(3n-6)$ and $|\Gamma(X)| \leq |X| + |S| = n-3 = 2|X| - \frac{n}{2}$. Similarly, if Y denotes the same set with one more vertex of $G - S$ added, then we also have $|Y| = \frac{1}{4}(3n-2)$ and $|\Gamma(Y)| \leq |Y| + |S| = n-1 = 2|Y| - \frac{n}{2}$.

Case 2. Suppose now that $|S| < \frac{1}{4}(n-6)$. Let h denote the number of vertices in all but the smallest component of $G - S$. Since there are $\geq |S| + 2$ components of $G - S$, containing between them $n - |S|$ vertices, we must have

$$h \geq \frac{|S| + 1}{|S| + 2} (n - |S|). \tag{6}$$

These h vertices can be adjacent to at most $h + |S| < n$ vertices; on the other hand, they are by hypothesis joined to at least $2h - \frac{n}{2}$ vertices. Thus

$$h \leq |S| + \frac{n}{2}. \tag{7}$$

From (6) and (7), eliminating h , we obtain

$$|S| \geq \frac{1}{4}(n-6),$$

a contradiction.

3. Extension to imperfect matchings

A related question is the following. Given a condition on a graph G which does not imply that G possesses a perfect matching, can we estimate how many disjoint edges can be found in G ? Corresponding to the two types of condition already studied, we have the following results for a graph with n vertices.

1. If each vertex is of degree $\geq cn$, $0 \leq c \leq \frac{1}{2}$, then we can find at least $\lceil cn \rceil$ disjoint edges.

2. If $\text{melt}(G) \geq c$, then there are at least

$$\frac{c}{c+1} n \text{ disjoint edges if } 0 < c \leq \frac{1}{2} \tag{8}$$

$$\left\lceil \frac{3c-2}{3c} n \right\rceil \text{ disjoint edges if } 1 < c \leq \frac{2}{3}. \tag{9}$$

Result 2 is due to Woodall (4), with (1) as the special case $c = \frac{4}{3}$. Result 1 is almost trivial (although best possible—consider a bipartite graph). For suppose that each vertex is of degree $\geq k$, and that $h < k$ disjoint edges have so far been found. If no two remaining vertices are joined by an edge, select any two of them, say v_1 and v_2 . Then it is easy to see that there must be a pair v_3, v_4 of vertices, joined by one of the edges already chosen, such that v_1 is joined to v_3 and v_2 to v_4 . With this new pairing we now have $h + 1$ disjoint edges, and the process can be repeated if $h + 1 < k$. We now state

Theorem 3. *Let G be a graph with n vertices. Suppose that*

(i) *each vertex is of degree $\geq dn$,*

(ii) $\text{melt}(G) \geq \frac{3-4d-3f}{2-2d}$,

where $4d + 3f \geq 1$, $2d + 3f \leq 1$, $d \geq 0$, $f \geq 0$. Then G possesses at least

$\left\lceil \frac{n}{2}(1-f) \right\rceil$ disjoint edges.

The special case $f = 0$ is Theorem 1, and the case $f = \frac{1}{3}(1-4d)$ is Woodall's result (9). The referee has suggested that it may be possible to deduce this result from an analogue to Theorem 2 in the same way as Theorem 1 was deduced from Theorem 2. However, we preserve here our original proof. Instead of Tutte's condition we use Berge's extension ((2); see also (4) for a simpler proof): for G to possess at least t disjoint edges, it is necessary and sufficient that $p(S) - |S| \leq n - 2t$ for all sets S of vertices of G . We shall in fact prove that, for all S ,

$$p(S) \leq |S| + nf + \frac{5}{3}$$

since this will imply that there are at least $\frac{n}{2}(1-f) - \frac{5}{6}$ and hence at least

$\left\lceil \frac{n}{2}(1-f) \right\rceil$ disjoint edges.

4. Proof of Theorem 3

In view of the above remarks, we may suppose that there exists a set S of vertices of G such that

$$p(S) > |S| + nf + \frac{5}{3} \tag{10}$$

and show that this leads to a contradiction.

Case 1. $|S| \geq dn$. Let m denote the number of 1-components in $G - S$. If $m = 0$,

$$n \geq |S| + 3p(S) > 4|S| + 3fn \geq (4d + 3f)n \geq n,$$

so we must have $m > 0$. Thus

$$n - m \geq |\Gamma(G - S)| \geq \frac{3-4d-3f}{2-2d}(n - |S|),$$

whence

$$m \leq \frac{3-4d-3f}{2-2d} |S| - \frac{1-2d-3f}{2-2d} n. \tag{11}$$

But we also have, from (3), ignoring the term $\frac{5}{3}$ in (10),

i.e.
$$\begin{aligned} n &> 4 |S| - 2m + 3nf, \\ m &> 2 |S| - \frac{1}{2}(1-3f)n. \end{aligned} \tag{12}$$

Eliminating m from (11) and (12), we obtain $|S| < dn$, a contradiction.

Case 2. $|S| < dn$. Here there can be no 1-components, so that each odd component contains at least

$$\max(3, dn - |S| + 1) \tag{13}$$

vertices. From now on we can assume that $4d + 3f > 1$.

Case 2(a). Suppose there is at least one 3-component. Then (13) yields

$$dn - |S| = \beta, \quad 0 < \beta \leq 2. \tag{14}$$

Then (3) and (10) give

$$n > 4 |S| + 3nf + 5 = (4d + 3f)n - 4\beta + 5$$

so that

$$n(4d + 3f - 1) < 4\beta - 5. \tag{15}$$

Considering on the other hand all but one of the odd components we have, from the definition of $\text{melt}(G)$,

$$n - 3 \geq \frac{3-4d-3f}{2-2d} (n - |S| - 3).$$

Substituting for $|S|$ from (14), this gives

$$n(1-d)(4d+3f-1) \geq (3-\beta)(4d+3f-1) - 6d + 2\beta > 2\beta - 6d.$$

Thus, by (15), we must have

$$2\beta - 6d < (1-d)(4\beta - 5)$$

whence

$$d > \frac{1}{3}.$$

It follows that

$$4d + 3f - 1 > \frac{1}{3}. \tag{16}$$

(15) and (16), with $\beta \leq 2$, now yield $n < 9$, and a contradiction easily follows.

Case 2 (b). Suppose now there is no 3-component. Here we shall show that $|S|$ is bounded. First of all, if $|S| < \frac{1}{2}dn$, then

$$n > |S| + (|S| + nf + \frac{5}{3})(\frac{1}{2}dn + 1)$$

so that

$$\frac{1}{2}dn |S| < |S|(2 + \frac{1}{2}dn) < n(1 - f - \frac{5}{6}d - \frac{1}{2}fdn). \tag{17}$$

If $dn < 4$ then $|S| < 1$ whereas, if $dn \geq 4$, then $\frac{1}{2}fdn \geq 2f$ and (17) yields $|S| < 6$. Secondly, if $\frac{1}{2}dn \leq |S| < dn$, then

$$n > |S| + (|S| + nf)(dn - |S| + 1)$$

whence

$$dn - |S| < \frac{n + nf}{|S| + nf} - 2 < \frac{1 + f}{\frac{1}{2}d + f} - 2 < 8.$$

Thus

$$8 + \frac{1}{2}n(1 - 5f) > 8 + |S| > dn,$$

$$40 > (6d + 5f - 1)n > \frac{1}{2}n,$$

so that $n < 80$. But

$$n \geq |S| + 5(|S| + 2)$$

so we must have $|S| \leq 11$.

Thus in any case, $|S| \leq 11$. It remains finally to consider each possible value of $|S|$ in turn. In each case we argue as follows.

Let h denote the number of vertices in all but one of the odd components of $G - S$. Then

$$|S| + h \geq \frac{3 - 4d - 3f}{2 - 2d} h$$

whence

$$|S| \geq \frac{1 - 2d - 3f}{2 - 2d} h.$$

Thus

$$|S| \geq \frac{1 - 2d - 3f}{2 - 2d} \cdot 5 \cdot (|S| + 1). \tag{18}$$

For any specific value of $|S|$, (18) gives a lower bound for d , and so for dn . For example, if $|S| = 5$, (18) yields

$$d \geq \frac{2}{5} - \frac{2}{5}f$$

and

$$dn \geq \frac{2}{5}n - \frac{2}{5}\theta$$

where $\theta = fn$. Having obtained this bound for dn , we obtain a contradiction by estimating n in two different ways. For we have

$$n > |S| + 5(|S| + \theta + 1)$$

and also

$$n > (|S| + \theta + 1)(dn - |S| + 1).$$

With $|S| = 5$, these give

$$n > 35 + 5\theta \tag{19}$$

and

$$n > (6 + \theta)(\frac{2}{5}n - \frac{2}{5}\theta - 4)$$

i.e.

$$n < \frac{9\theta^2 + 74\theta + 120}{2\theta + 7}. \tag{20}$$

(19) and (20) contradict one another. The theorem is now proved.

REFERENCES

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