A NOTE ON HOMOGENEOUS DENDRITES

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1. In graph-theoretic terms a homogeneous p-dendrite, $p \ge 2$, is defined as a finite singly-rooted tree in which the root has valency 1 while every other vertex has valency 1 or p. More descriptively, a homogeneous p-dendrite may be imagined to start from its root as the main, or 0th order, branch which proceeds to the first-order branch point where it gives rise to p first-order branches. Each of these either terminates at its other end (which is a second-order branch point) or it splits there again into p branches (which are of third order), and so on. The order of the dendrite is the highest order of a branch present in it. For completeness, a 0-th order dendrite is also allowed, this consists of the 0-th order branch alone.

Alternatively, if we consider a development in time rather than a structure in space, a homogeneous p-dendrite represents a history in which a single individual fissions into p identical individuals each of which either dies without descendants or else, fissions into p new indistinguishable individuals again.

We shall be interested here in the number f (n) of (topologically) distinct n-th order p-dendrites. Our interest is motivated partly by biological and physical considerations relative to certain simple branching processes (number of various family-histories, number of distinct dendrites of a neuron, particle-showers, etc.) and partly by pure combinatorics.

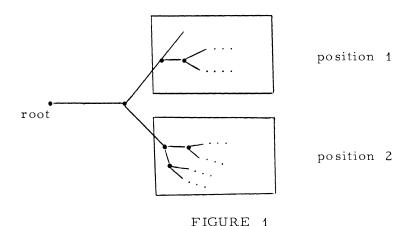
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2. Here we determine the number $f_2(n) = f(n)$ of binary dendrites. To begin with, we have

(1)
$$f(0) = 1, f(1) = 1.$$

Let $n \ge 0$ and consider an (n+1)-st order binary dendrite. There is here one first-order branch point, as shown in Figure 1, and this is followed by two structures one of which is an n-th order dendrite (position 1) and the other one an m-th order dendrite (position 2), with m < n.



Suppose first that m < n. Then position 1 can be filled by any one of the f(n) distinct n-th order dendrites and, independently, position 2 by any one of the f(m) m-th order ones. Hence the total number of (n+1)-st order dendrites with m < n is

(2)
$$N_{1} = f(n) \sum_{i=0}^{n-1} f(i)$$

When m = n, the n-th order dendrites in positions 1 and 2 can be identical, and this can occur in f(n) ways, or they can be distinct, which can occur in

$$f(n)[f(n) - 1]/2$$

ways. Therefore the total number of (n+1)-st order dendrites with m = n is

$$N_2 = f(n)[f(n) + 1]/2.$$

Adding N_1 and N_2 we get the total number f(n+1) of distinct (n+1)-st order dendrites:

(3)
$$f(n+1) = f(n) \begin{bmatrix} n-1 \\ \sum_{i=0}^{n-1} f(i) + (1+f(n))/2 \end{bmatrix},$$

which may be written as

(4)
$$\frac{f(n+1)}{f(n)} = \frac{1 + f(n)}{2} + \sum_{i=0}^{n-1} f(i).$$

Copying this equation with n replaced by n-1, subtracting from (4), and re-arranging, we get a nonlinear second-order recursion for f:

(5)
$$f(n+1) = f(n) \left[\frac{f(n)}{f(n-1)} + \frac{f(n) + f(n-1)}{2} \right].$$

From this and (1) we compute successively

$$f(2) = 2$$
, $f(3) = 7$, $f(4) = 56$, $f(5) = 2212$, $f(6) = 2595782$

and so on. It would be useful to have an explicit formula for f(n) but this does not appear to be easy to get. Some rough bounds on f(n) can be obtained as follows. By (5) we have

(6)
$$f^{2}(n_{0})/2 \le f(n_{0}+1) \le 2f^{2}(n_{0});$$

therefore

$$[f^{2}(n_{o})/2]^{2}/2 \le f(n_{o}+2) \le 2[2f^{2}(n_{o})]^{2}$$

and generally for arbitrary n and any fixed n

(7)
$$[f^{2^{n}}(n_{0})]/2^{2^{n}-1} \le f(n+n_{0}) \le 2^{2^{n}-1}[f^{2^{n}}(n_{0})].$$

3. Throughout this section we assume that p=3 and we get a formula analogous to (5) for the case of ternary splitting. Let $f_3(n) = f(n)$, put $n \ge 2$ and consider an (n+1)-st order ternary dendrite. In analogy to Figure 1 we have now three positions to be filled by ternary dendrites of orders n, m, and k, with $0 \le k \le m \le n$. Suppose first that k < m < n, then any k-th order, m-th order, and n-th order dendrites can fill, independently, their respective positions and so the number N_1 of ternary dendrites of (n+1)-st order, with k < m < n. is

(8)
$$N_{1} = f(n) \sum_{m=1}^{n-1} \sum_{k=0}^{m-1} f(k)f(m).$$

When k < m = n the corresponding number is

(9)
$$N_2 = f(n) \left[\frac{f(n) + 1}{2} \right] \sum_{k=0}^{n-1} f(k).$$

When k=m=n, there are three cases to consider because among the n-th order dendrites filling the three positions there may be one, two or three distinct ones. The total number N_3 is here

(10)
$$N_3 = f(n) + f(n)[f(n) - 1] + f(n)[f(n) - 1][f(n) - 2]/6$$

= $f(n)[f^2(n) + 3f(n) + 2]/6$.

Finally, when k = m < n, the contribution to the total is

(11)
$$N_4 = \frac{f(n)}{2} \sum_{k=0}^{n-1} f(k)[f(k) + 1].$$

Adding the numbers N_1 , N_2 , N_3 , N_4 from the equations (8), (9), (10), (11) we get

$$f(n+1) = f(n) \sum_{m=1}^{n-1} \sum_{k=0}^{m-1} f(m)f(k) + \frac{f(n)[f(n)+1]}{2} \sum_{k=0}^{n-1} f(k) + \frac{f(n)}{2} \sum_{k=0}^{m-1} f(k)[f(k)+1] + \frac{f(n)}{6} [f^{2}(n)+3f(n)+2].$$

Therefore

(13)

$$\frac{f(n+1)}{f(n)} - \frac{f^{2}(n) + 3f(n) + 2}{6} = \sum_{m=1}^{n-1} \sum_{k=0}^{m-1} f(m)f(k) + \frac{f(n) + 1}{2} \sum_{k=0}^{n-1} f(k)$$

$$+ \frac{1}{2} \sum_{k=0}^{n-1} f(k)[f(k) + 1].$$

Denote the left-hand side of (13) by F(n); taking first differences, we get

$$F(n) - F(n-1) = f(n-1) \sum_{k=0}^{n-2} f(k) + \frac{f(n)+1}{2} f(n-1) + \frac{f(n)-f(n-1)}{2} \sum_{k=0}^{n-2} f(k) + \frac{1}{2} f(n-1)[f(n-1)+1]$$
so that

(14)

$$F(n) - F(n-1) - \frac{f(n-1)}{2} [f(n) + f(n-1)] - f(n-1) = \frac{f(n) + f(n-1)}{2} \sum_{k=0}^{n-2} f(k).$$

Denote the left-hand side of (14) by G(n) and put

$$H(n) = \frac{2G(n)}{f(n) + f(n-1)}$$

so that (14) is now simply

$$H(n) = \sum_{k=0}^{n-2} f(k);$$

taking first differences again, we eliminate all the sums and get

$$H(n) - H(n-1) = f(n-2).$$

Substituting successively for H. G, F, we get after some tedious algebra

(15)

$$\begin{split} f(n+1) &= f(n)[f(n) + f(n-1)][f(n) + f(n-1) + f(n-2)]/6 \\ &+ f(n) \frac{f(n) + f(n-1)}{f(n-1) + f(n-2)} \left[\frac{f(n)}{f(n-1)} - \frac{f(n-1)}{f(n-2)} \right] + f^2(n)/f(n-1). \end{split}$$

Direct inspection shows that f(0) = 1, f(1) = 1, f(2) = 3, now the recursion formula (15) yields

$$f(3) = 31$$
, $f(4) = 8401$, $f(5) = 100130704103$ etc.

4. It is possible to obtain in the same way successive formulas, analogous to (5) and (15), for $f_4(n)$, $f_5(n)$, etc.

However their complexity grows very rapidly, and a recursion formula valid for a general f (n) appears to be difficult to get.

We shall obtain instead the general analogue of (3) and (12).

Consider an (n+1)-st order p-dendrite. Referring to Figure 1. we have here p positions to fill instead of two, and we suppose that the j-th position contains a p-dendrite of order k_i . To meet the enumerative conditions we must have

(16)
$$0 \le k_1 \le k_2 \le \ldots \le k_p = n.$$

It is important to know where the strict inequality occurs between k_i and $k_{i+1},\ i=1,\ldots,p-1.$ There are 2^{p-1} sequences of p-1 signs each of which is "<" or "="; any such sequence will be denoted by r and called an ordering, and the set of all 2^{p-1} orderings will be denoted by R. Once an ordering r ε R

is given, the monotonicity properties of (16) are known; further, irrespective of the values of the indices k_i two different orderings will lead to different (n+1)-st order dendrites. Let $r \in \mathbb{R}$, if in r we find a sequence such as

$$k_1 = k_2 = \dots = k_m < \dots, \dots < k_i = k_{i+1} = \dots = k_{i+m-1} < \dots,$$

$$or \dots < k_{p-m+1} = k_{p-m+2} = \dots = k_p$$

we call it a step of length m. In particular,

$$k_1 < \ldots, \ldots < k_i < \ldots, \ldots < k_p$$

are steps of length 1. Let g = g(r) be the total number of steps in r and let $m_j = m_j(r)$ be the length of the j-th consecutive one (j = 1, 2, ..., g(r)) so that

$$1 \le g(r) \le p, \quad \sum_{j=1}^{g(r)} m_j(r) = p.$$

Consider now the j-th step, of length $m_j(r)$; this corresponds to filling $m_j(r)$ positions with p-dendrites of the same order, say s_j . By the enumeration conditions of the problem we deal here with combinations in which repetitions are allowed. and there are $f_p(s_j)$ possibilities of filling each position. Therefore the $m_j(r)$ positions can be filled in

$$\begin{pmatrix}
f_{p}(s_{j}) + m_{j}(r) - 1 \\
m_{j}(r)
\end{pmatrix}$$

ways. Hence the number of ways in which all the positions can be filled, once the ordering r as well as the values of the indices s; are fixed, is

$$\frac{g(\mathbf{r})}{\Pi} \left(f_{\mathbf{p}}(\mathbf{s}_{j}) + m_{j}(\mathbf{r}) - 1 \right)$$

$$\frac{g(\mathbf{r})}{\Pi} \left(f_{\mathbf{p}}(\mathbf{s}_{j}) + m_{j}(\mathbf{r}) - 1 \right)$$

Allowing for suitable variation of indices s corresponding to the same ordering r, we find that the total number of ways of filling all the positions for a fixed ordering r is

Summing over all the 2^{p-1} orderings to get the grand total number of ways of filling all the positions we get finally

(17)

$$f_{p}(n+1) = \sum_{\mathbf{r} \in R} \begin{bmatrix} n-1 & s_{4}^{-1} & s_{3}^{-1} & s_{2}^{-1} g(\mathbf{r}) & f_{p}(s_{j}) + m_{j}(\mathbf{r}) - 1 \\ \sum & \dots & \sum & \sum & \prod & m_{j}(\mathbf{r}) \\ s_{g(\mathbf{r})-1} = g(\mathbf{r}) - 2 & s_{3} = 2 & s_{2} = 1 & s_{1} = 0 & j = 1 \end{bmatrix} \begin{pmatrix} f_{p}(s_{j}) + m_{j}(\mathbf{r}) - 1 \\ m_{j}(\mathbf{r}) \end{pmatrix}$$

which generalizes (3) and (12) to arbitrary p.

Of the 2^{p-1} terms in the square brackets there is exactly one, namely

$$\begin{pmatrix} f_p(n) + p - 1 \\ p \end{pmatrix}$$

containing no summation; this corresponds to having all p positions filled with maximal (n-th order) dendrites. Therefore

(18)

$$f_{p}(n+1) - \begin{pmatrix} f_{p}(n) + p - 1 \\ p \end{pmatrix} = \sum_{\substack{p-1 = p-2 \\ s_{p-1} = p-2}}^{n-1} \sum_{\substack{s_{4}^{-1} \\ s_{3}^{-2} = 2}}^{s_{4}^{-1} \cdot s_{3}} \sum_{\substack{s_{2}^{-1} \\ s_{1}^{-2} = 0}}^{n-1} g(r) \\ f_{p}(s_{j}) + Q_{p-2}(s_{j}) + Q_{p-2}(s_$$

where the first term on the right corresponds to all p positions having dendrites of different orders (assuming that n is large enough) and has p-1 summations, while Q_{p-2} is the sum of all the other terms, each of which has $\leq p-2$ summations. Denote the left-hand side of (18) by F(n); taking first differences one

finds that F(n) - F(n-1) is of the form

(19)
$$f_{p}(n-1)S_{1} + T_{p-3} + T_{p-4} + \ldots + T_{1} + T_{0}$$

where S_1 is a single term with p-2 summations, and T_i is a sum of terms with i summations. One repeats now the same number-of-summations reduction procedure by taking the first difference of

$$G(n) = [F(n) - F(n-1) - T_0]/f_p(n-1)$$

to get an expression similar to (19):

$$G(n) - G(n-1) = \varphi[f_p(n), f_p(n-1)]S'_1 + T'_{p-4} + ... + T'_0$$

where $\,\varphi\,$ is a rational function. The whole process is carried out p-1 times and one ends up with a nonlinear p-step recurrence relation

$$f_{p}(n+1) = R[f_{p}(n), f_{p}(n-1), ..., f_{p}(n-p+1)]$$

where R is a rational function with integer coefficients.

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